Easy, Effective, Efficient: GPU Programming in Python with PyOpenCL and PyCUDA

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Courant Institute of Mathematical Sciences
New York University

Simula PyOpenCL Workshop
Lecture 1 · August 23, 2011

Course Outline

Morning Session: Intro

- Python, numpy, GPUs
- OpenCL
- Basic PyOpenCL
- Tour of PyOpenCL Runtime
- Advanced PyOpenCL usage
- OpenCL device language
- PyOpenCL: Built-in tools
- CL Implementation Notes

Lunch Lab

- Python, numpy
- Basic PyOpenCL

Afternoon Session: Advanced

- Behind the scenes
- RTCG: How and Why, Templating
- Automated Tuning
- mpi4py and PyOpenCL
- Interfacing Python with Fortran and C/C++
- A brief look at PyCUDA

Afternoon Lab

- Continue on Lab 1
- Advanced PyOpenCL

Outline

Outline

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Outline

Python in 4 Minutes

Literals `1234`, `1234.`, `0xabc` `"a string"` `"""a multi-line string"""` `["a", "list"]` `("a", "tuple", 17)` `{"a": 17, "dictionary": 19}`

Flow Control

```
if True and a == 10:
   print "?" # a comment
while 0 <= x < 17 :
   pass # break, continue
for i in [0, 1, 2]:
   raise Exception("!")
```

Functions, Classes

```
def my_function(x):
   return 17*x
class MyClass:
   def __init__(self, x):
      self.x = x
```

Program Semantics

```
a = [1,2,4]
b = a
b.append(17)
print a
# [1, 2, 4, 17]
```

Python in 4 Minutes

Literals `1234`, `1234.`, `0xabc`
```
"a string" """a multi-line
string""" ["a", "l
("a", "tuple", 17)
{"a": 17, "diction
```

Functions, Classes
```
def my_function(x):
    return 17*x
```

Flow Control
```
if True and a == 1
    print "?" # a c
while 0 <= x < 17
    pass # break, c
for i in [0, 1, 2]
    raise Exceptio
```

`http://docs.python.org`

More stuff:

- Python 2 vs Python 3
- 'Batteries included'
- The package index
- Cython, Jython, IronPython, PyPy
- Interactive console, IPython, PuDB, Virtualenv, Pip, Spyder, PEP 8

Numpy in 4 Minutes

Creating/Modifying Arrays

```
import numpy as np
x = np.array([[1,2],[4,5]])
print x.shape # (2,2)

y = np.zeros((20000, 3),
  dtype=np.float64)
z = np.empty((20000, 3))
u = np.ones((30, 40))
v = np.linspace(1, 5,20,
  endpoint=False)

# also: mgrid, eye, arange
+, -, *, +=, np.dot
```

Indexing Arrays

```
a = x[:, 1] # a 'view'
a[:, :] = 17
y = 17
x[3:-3:-1, :] = 17
x[x == 19] = 17
```

Broadcasting

```
y[:, :] = 17
y[:, :] = [0, 1, 2]
w = np.array([0, 1, 2]) \
  [:, np.newaxis] * [0, 1, 2]
```

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Numpy in 4 Minutes

Creating/Modifying Arrays

```
import numpy as np
x = np.array([[1,2
print x.shape # (2

y = np.zeros((2000
  dtype=np.float64
z = np.empty((2000
u = np.ones((30, 4
v = np.linspace(1,
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# also: mgrid, eye
+, -, *, +=, np.dot
```

Indexing Arrays

```
a = x[:, 1] # a 'view'
```

2]

`http://docs.scipy.org`

More stuff:

- 'ufuncs' `sin,exp,...`
- Linear Algebra, FFT, ..., SciPy
- Structured/masked arrays
- 'Fancy' Indexing
- Matplotlib, MayaVi2
- C API
- Google 'Numpy Medkit'

Questions?

?

Outline

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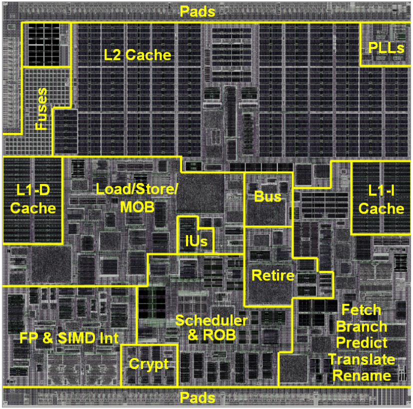
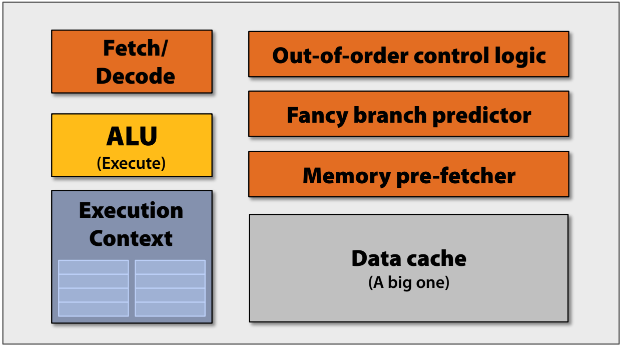
CPU Chip Real Estate



Die floorplan: VIA Isaiah (2008).
65 nm, 4 SP ops at a time, 1 MiB L2.

"CPU-style" Cores



Credit: Kayvon Fatahalian (Stanford)

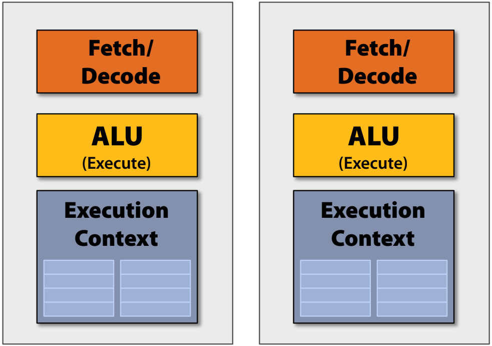
Slimming down

Idea #1:

Remove components that help a single instruction stream run fast

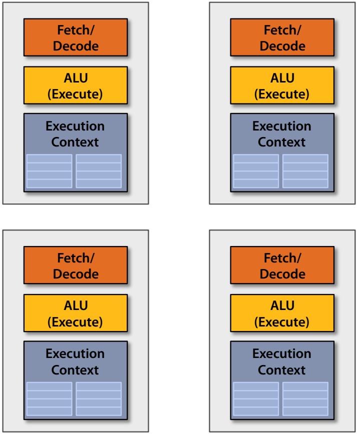
Credit: Kayvon Fatahalian (Stanford)

More Space: Double the Number of Cores



Credit: Kayvon Fatahalian (Stanford)

. . . again



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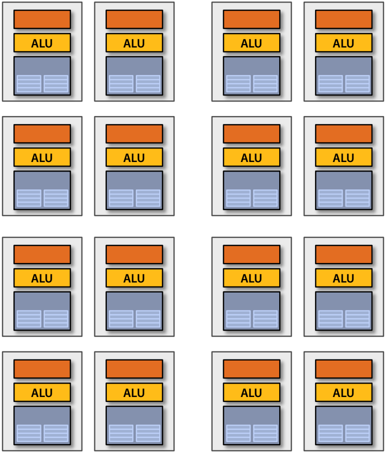

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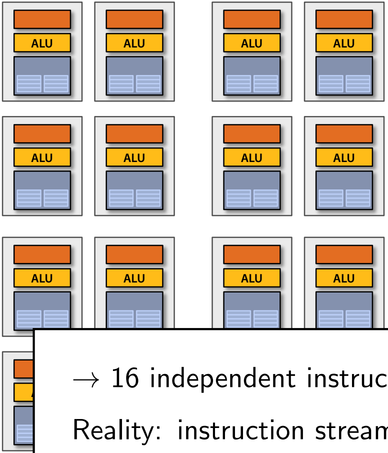
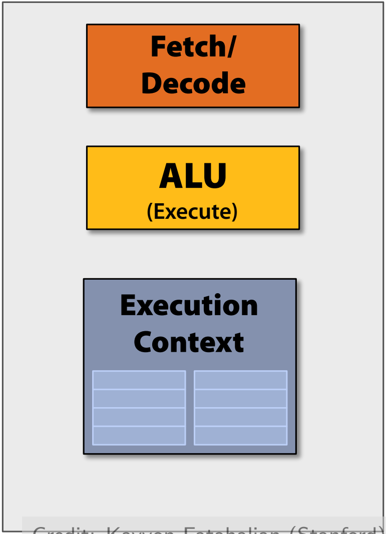


→ 16 independent instruction streams

Reality: instruction streams not actually very different/independent

Credit: Kayvon Fatahalian (

Saving Yet More Space



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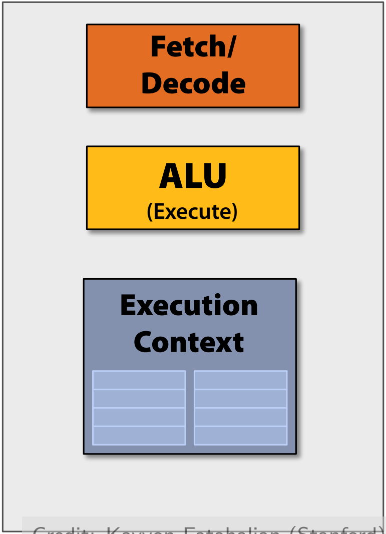

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Idea #2

Amortize cost/complexity of managing an instruction stream across many ALUs

→ **SIMD**

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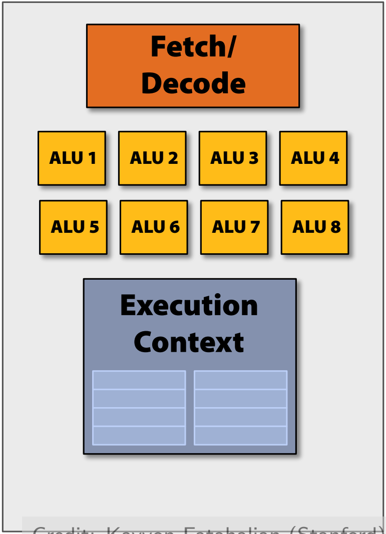

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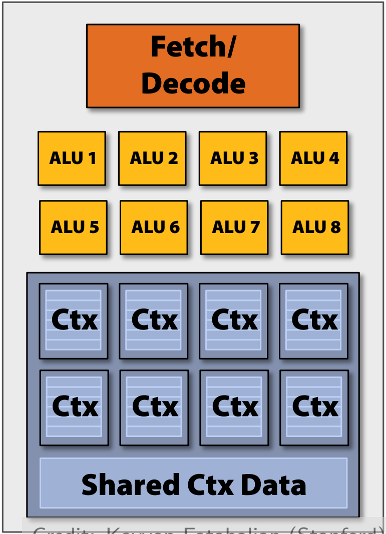

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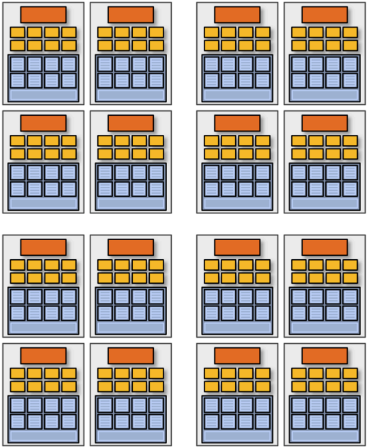
$\rightarrow$ **SIMD**

Gratuitous Amounts of Parallelism!



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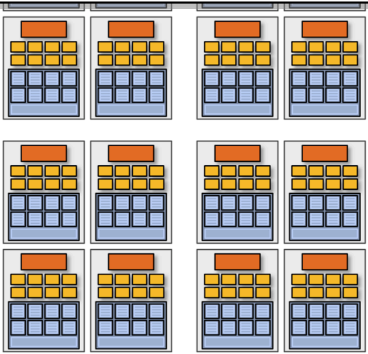
Gratuitous Amounts of Parallelism!

Example:

128 instruction streams in parallel
16 independent groups of 8 synchronized streams



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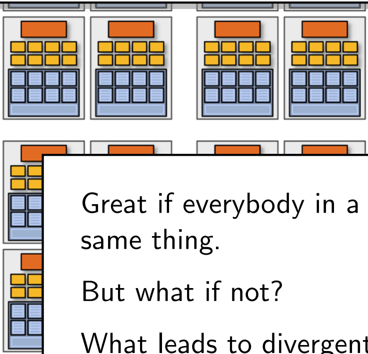


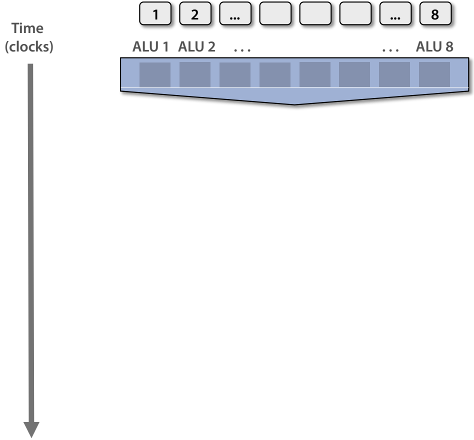
Great if everybody in a group does the same thing.

But what if not?

What leads to divergent instruction streams?

Credit: Kayvon Fatahalian (

Branches




```
<unconditional
shader code>

if (x > 0) {
    y = pow(x, exp);
    y *= Ks;
    refl = y + Ka;
} else {
    x = 0;
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<resume unconditional
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```

Credit: Kayvon Fatahalian (Stanford)

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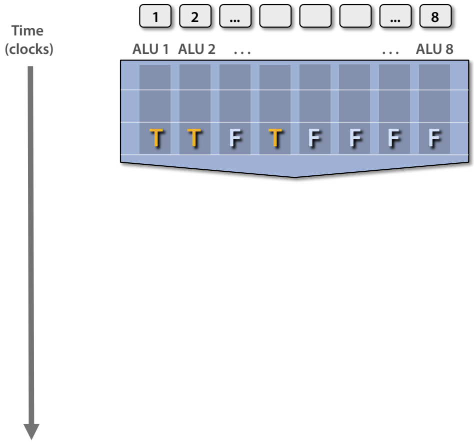


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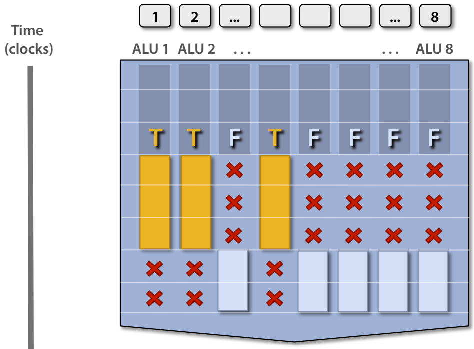


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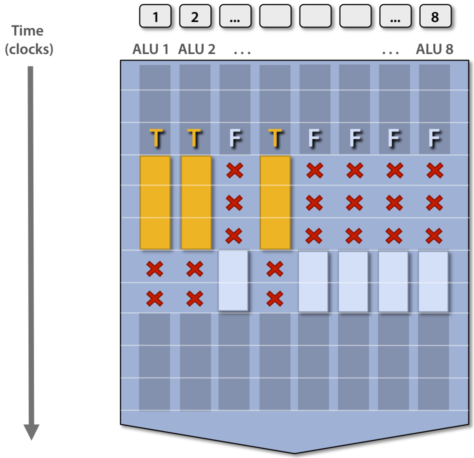
Not all ALUs do useful work!
Worst case: 1/8 performance

Credit: Kayvon Fatahalian (Stanford)

Branches

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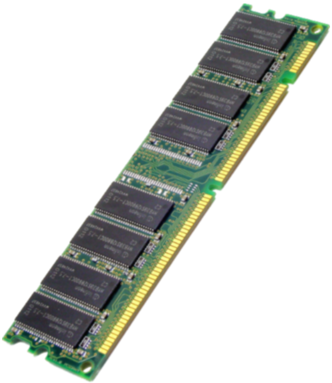
Remaining Problem: Slow Memory

Problem

Memory still has very high latency. . .
. . . but we've removed most of the hardware that helps us deal with that.



We've removed

- caches
- branch prediction
- out-of-order execution

So what now?

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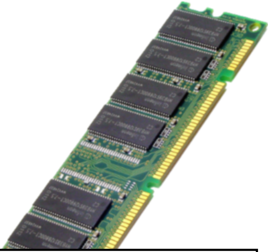

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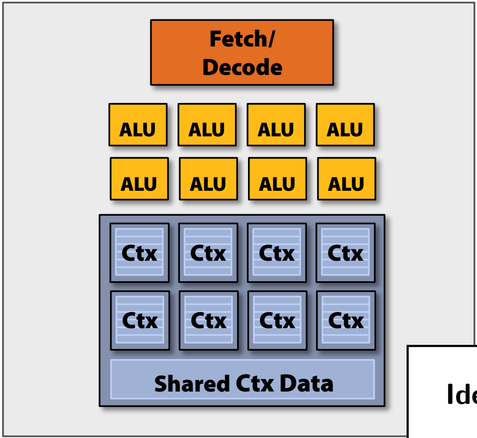
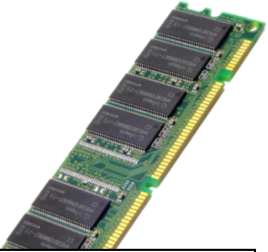
So what now?

Idea #3

	Even more parallelism
$+$	Some extra memory
$=$	A solution!





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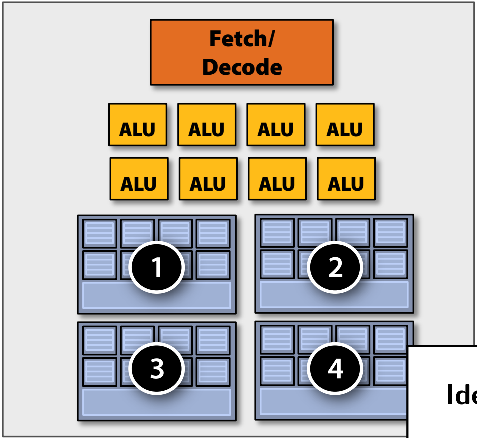

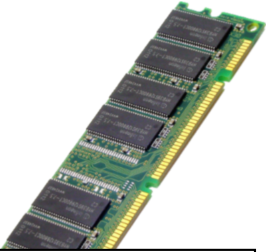

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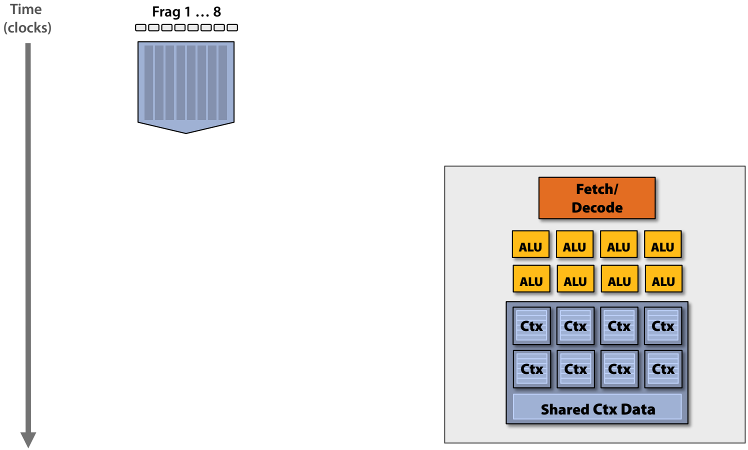
Hiding Memory Latency



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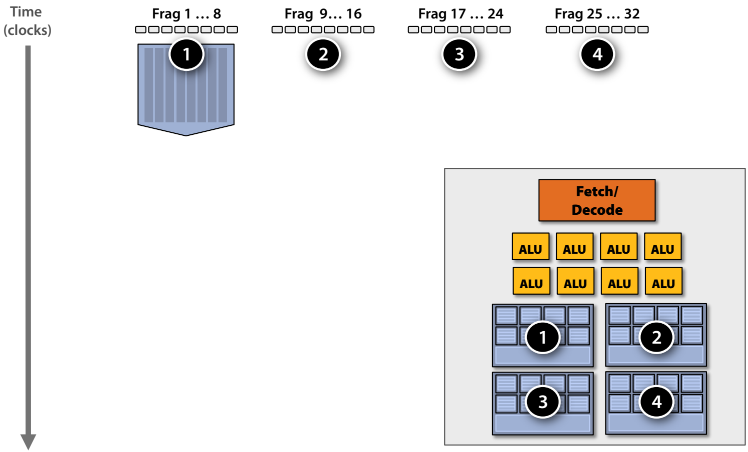

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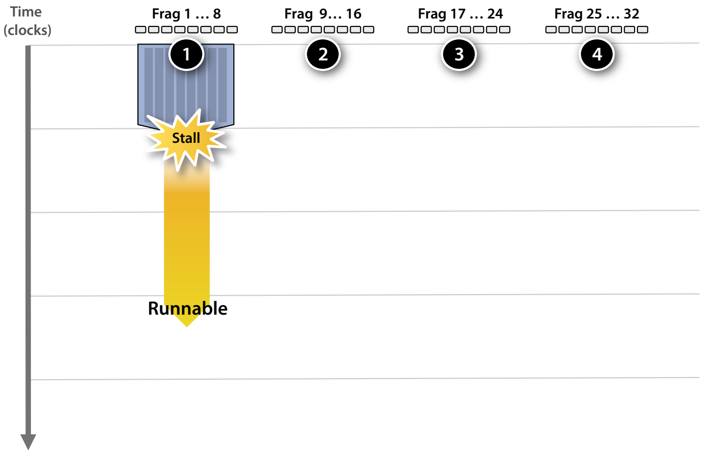

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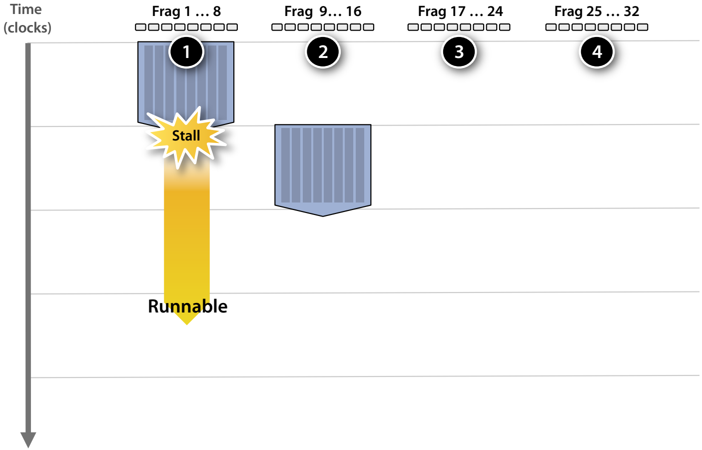

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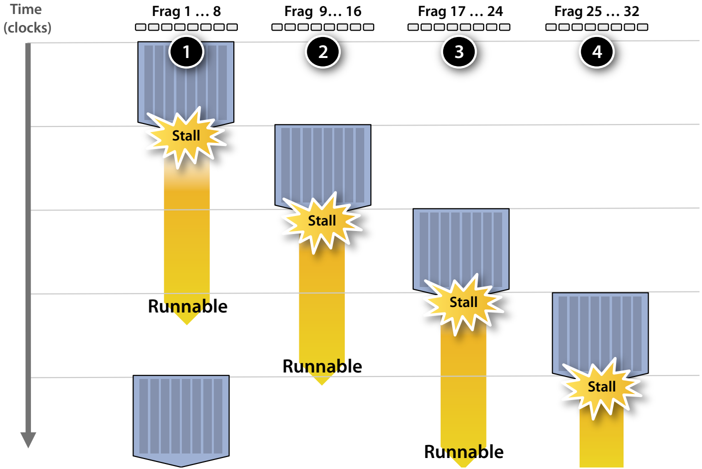

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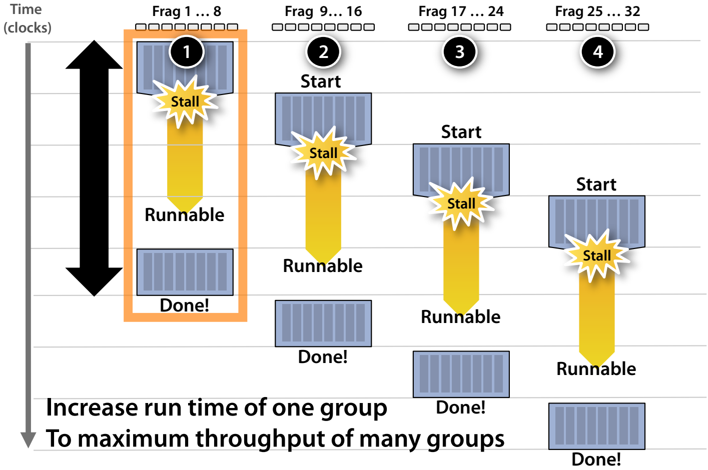


Credit: Kayvon Fatahalian (Stanford)

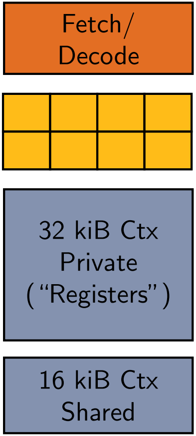
GPU Architecture Summary

Core Ideas:

1. Many slimmed down cores
   $\rightarrow$ lots of parallelism
2. More ALUs, Fewer Control Units
3. Avoid memory stalls by interleaving execution of SIMD groups ("warps")

Credit: Kayvon Fatahalian (Stanford)

Connection: Hardware ↔ Programming Model



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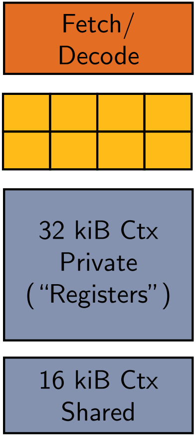

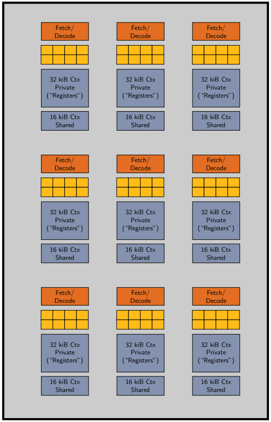

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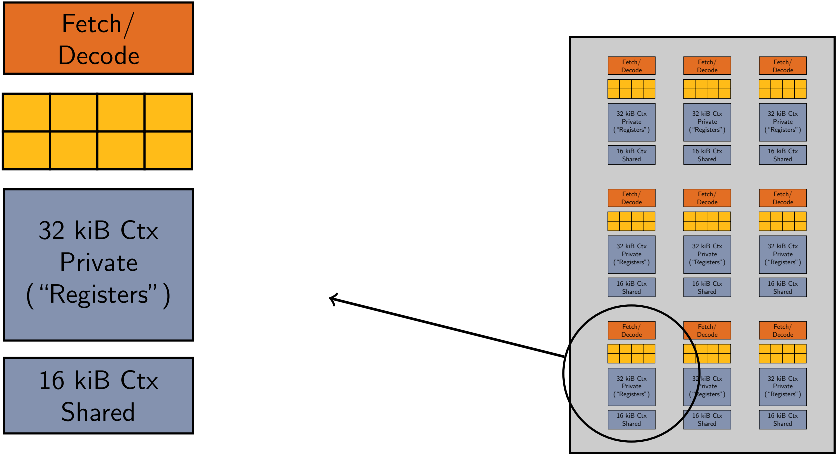

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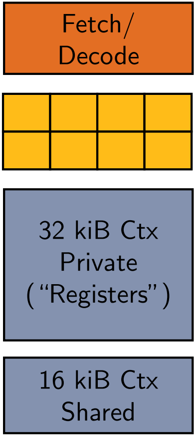

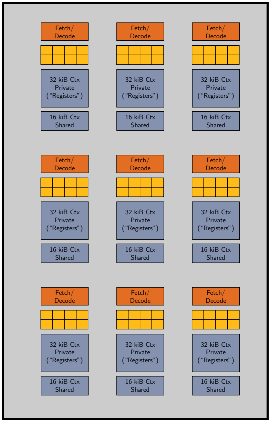

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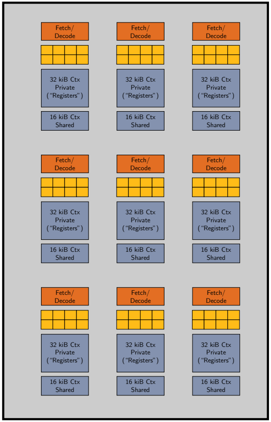
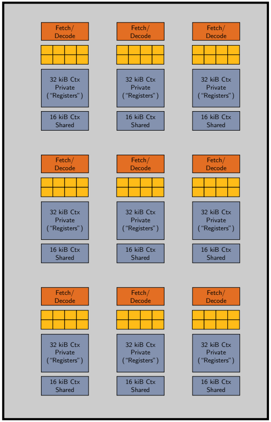


Connection: Hardware ↔ Programming Model

Who cares how many cores?

Idea:

- Program as if there were "infinitely" many cores
- Program as if there were "infinitely" many ALUs per core



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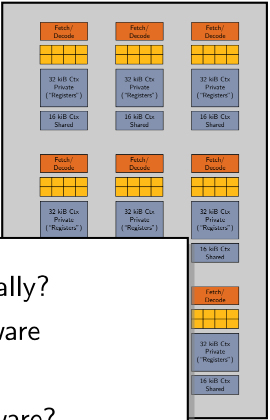

Consider: Which is easy to do automatically?

- Parallel program → sequential hardware

or

- Sequential program → parallel hardware?

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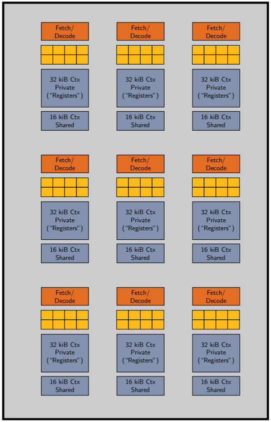

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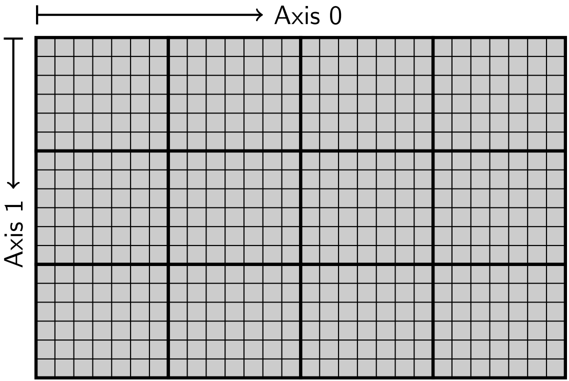

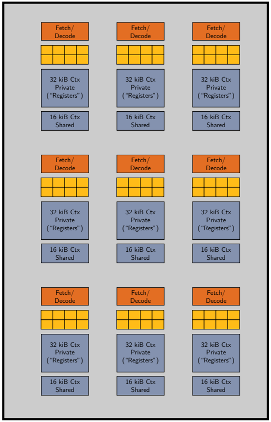

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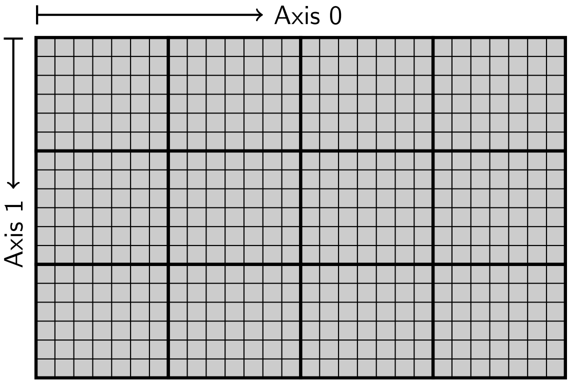

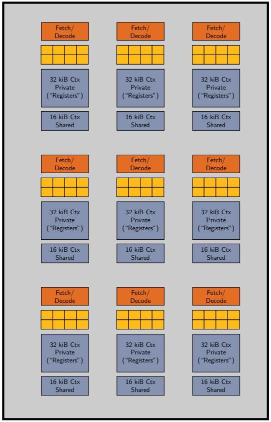

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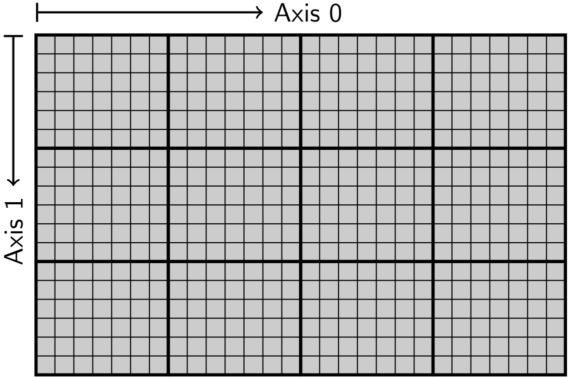
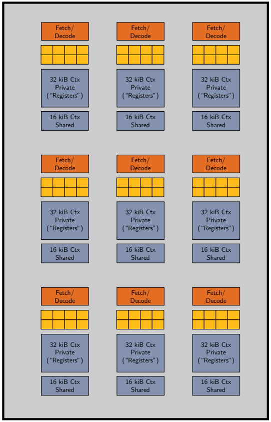


Software representation



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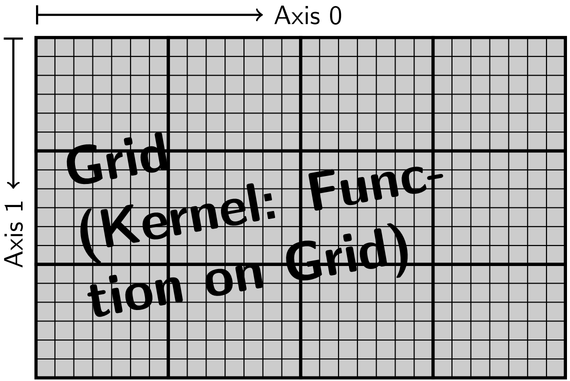

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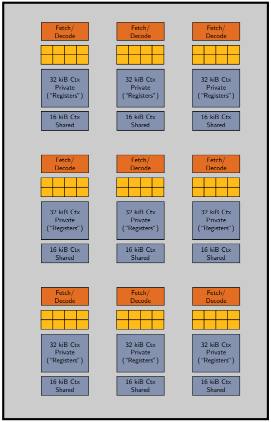

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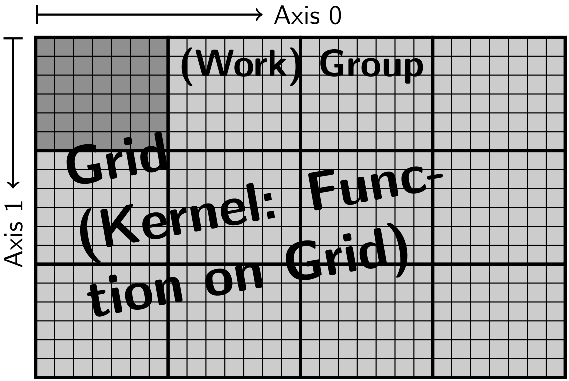

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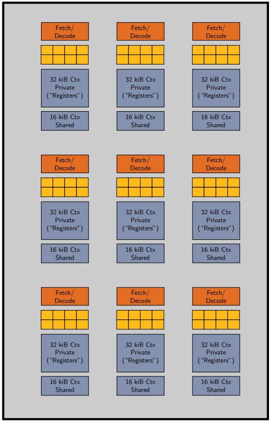

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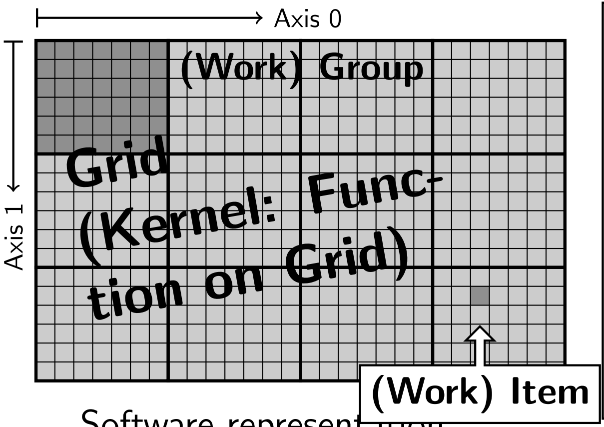
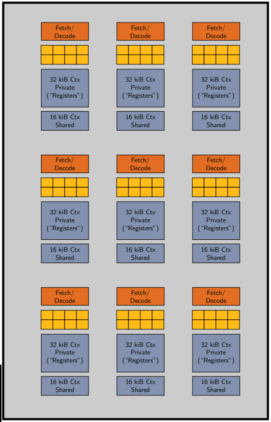
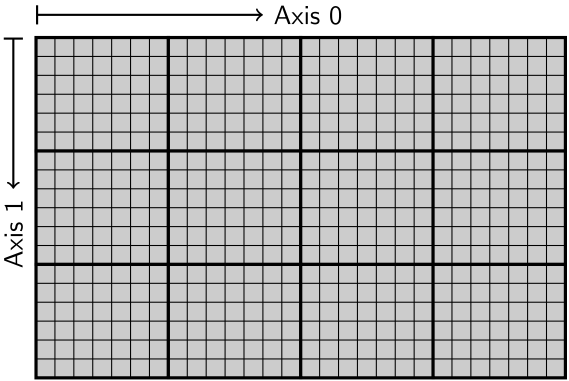


Software representation

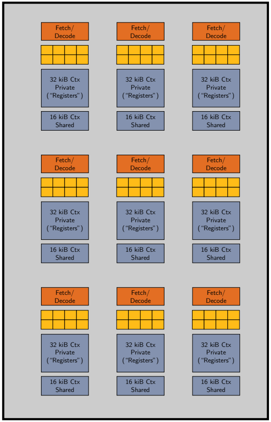


Hardware

Connection: Hardware ↔ Programming Model

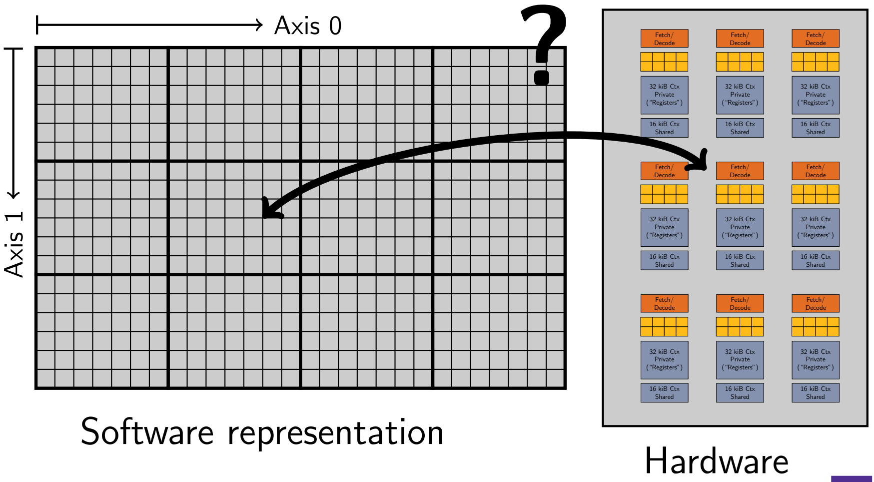


Software representation



Hardware

Connection: Hardware ↔ Programming Model



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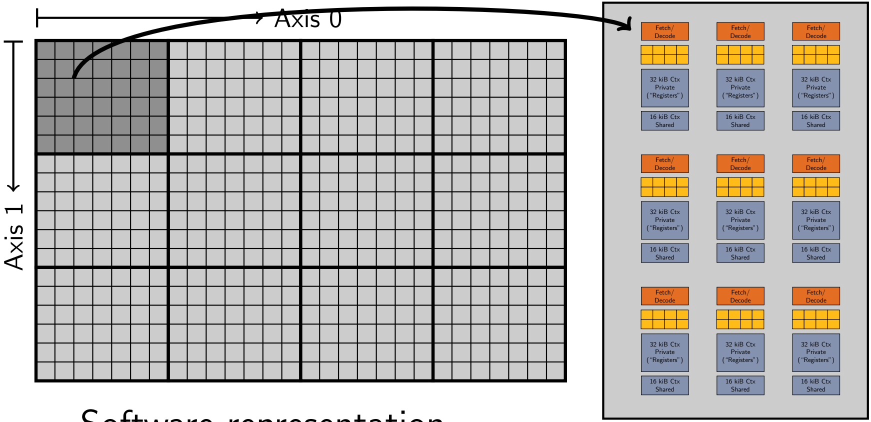

Software representation

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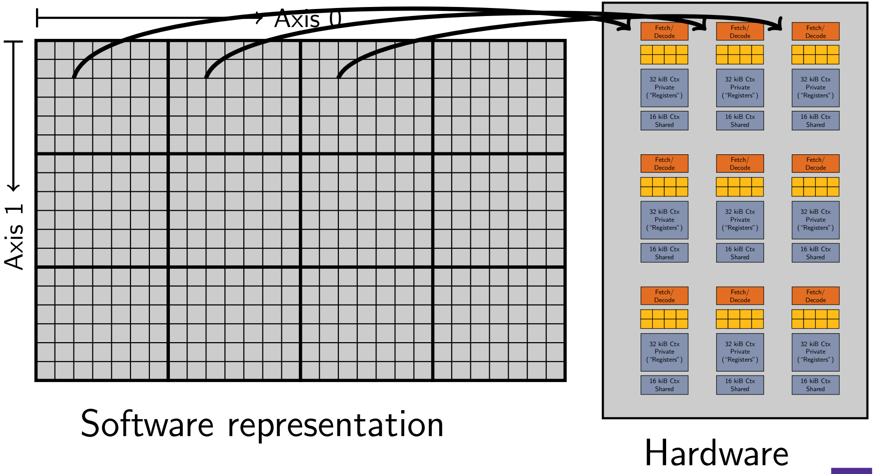

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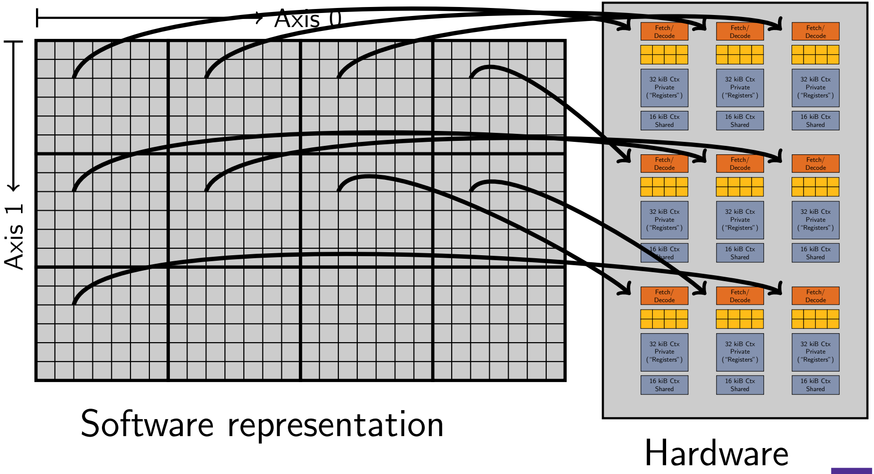

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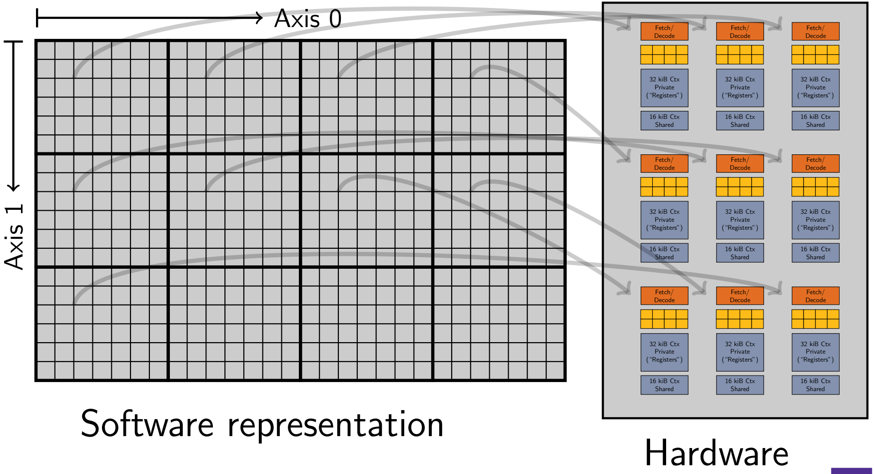

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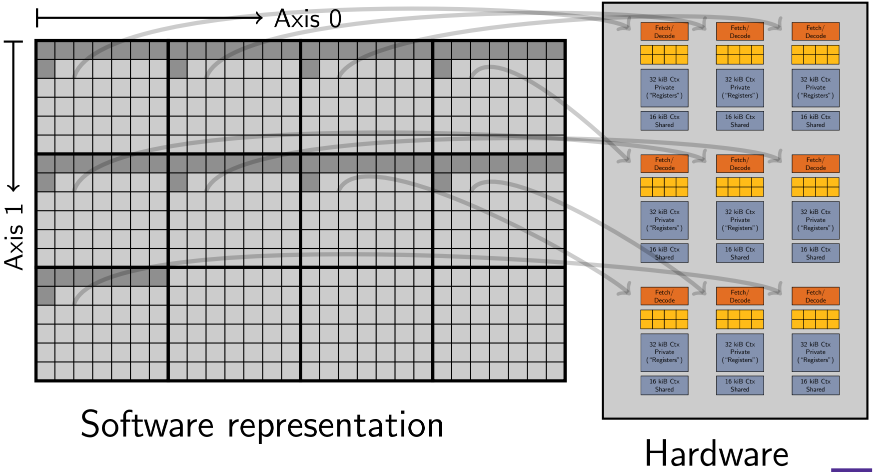

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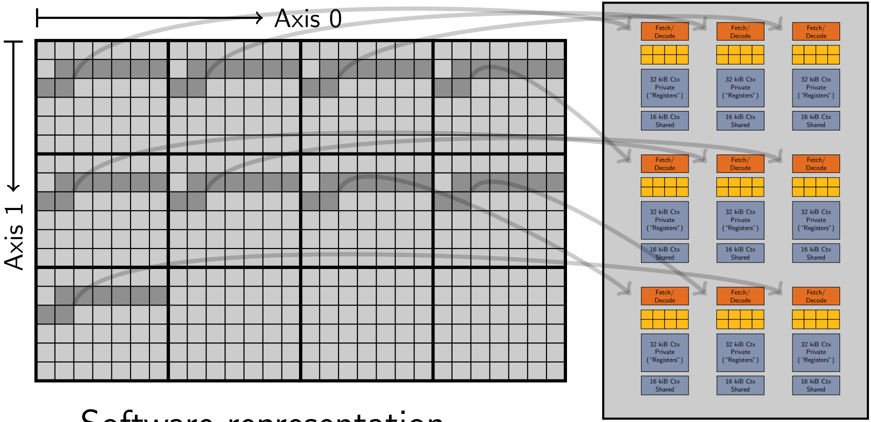
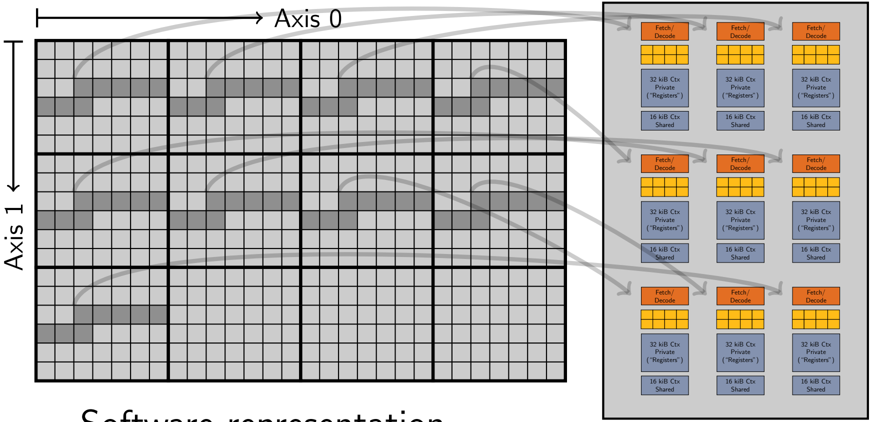


Software representation

Hardware

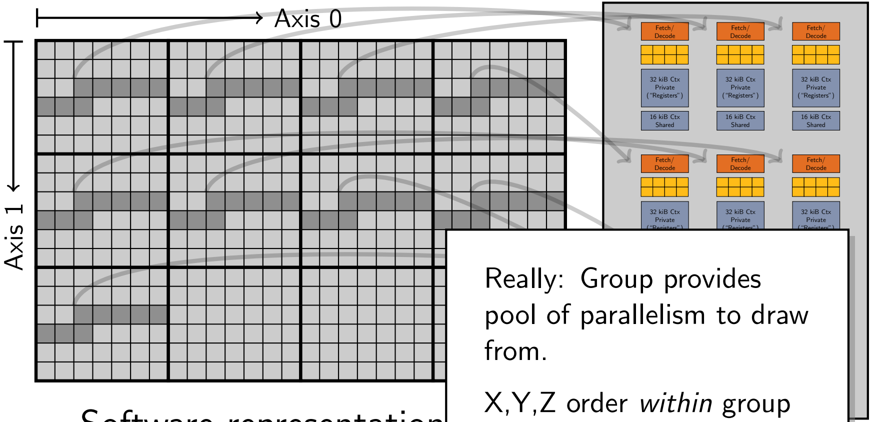
Connection: Hardware ↔ Programming Model



Software representation

Hardware

Connection: Hardware ↔ Programming Model

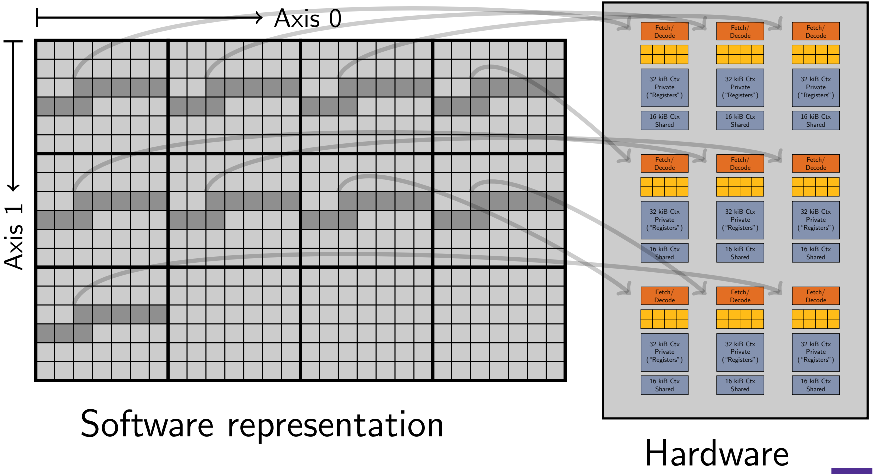


Software representation

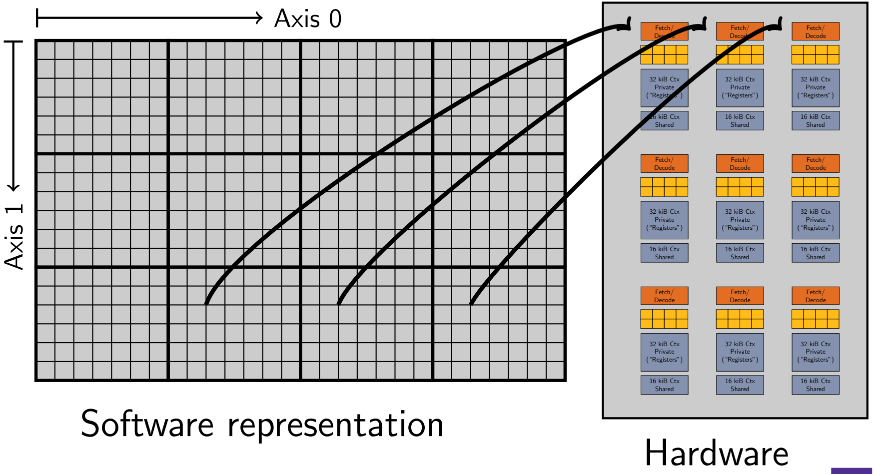
Really: Group provides pool of parallelism to draw from.

X,Y,Z order *within* group matters. (Not *among* groups, though.)

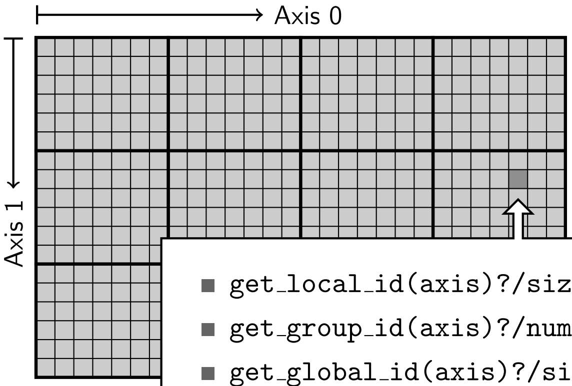
Connection: Hardware ↔ Programming Model

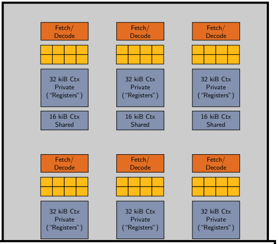


Connection: Hardware ↔ Programming Model



Connection: Hardware ↔ Programming Model





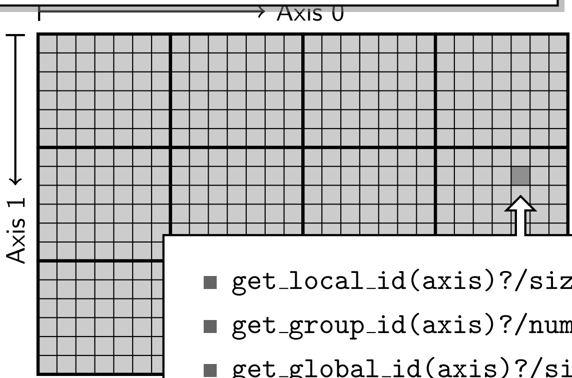
- `get_local_id(axis)?/size(axis)?`
- `get_group_id(axis)?/num_groups(axis)?`
- `get_global_id(axis)?/size(axis)?`

`axis=0,1,2,...`

Hardware

Co… g Model

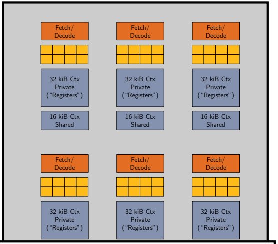
Grids can be 1,2,3-dimensional.



- `get_local_id(axis)?/size(axis)?`
- `get_group_id(axis)?/num_groups(axis)?`
- `get_global_id(axis)?/size(axis)?`

`axis=0,1,2,...`

Softw…



Hardware

GPU architecture: Overview

Now know about basic execution model.

Observe: Same model also applies to multi-core CPUs!

$\rightarrow$ the "OpenCL" execution model

Will learn more about GPUs later. In particular:

- Memory access
- Device Management
- Synchronization

Note: CPUs have a very different memory system.

Outline

What is OpenCL?

OpenCL (Open Computing Language) is an open, royalty-free standard for general purpose parallel programming across CPUs, GPUs and other processors. [OpenCL 1.1 spec]

- Device-neutral (Nv GPU, AMD GPU, Intel/AMD CPU)
- Vendor-neutral
- Comes with RTCG

Defines:

- Host-side programming interface (library)
- Device-side programming language (!)

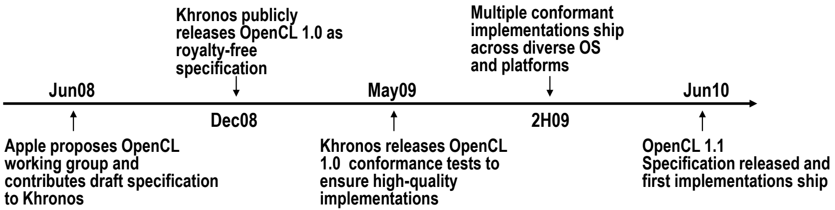
Who?

- **Diverse industry participation**
  - Processor vendors, system OEMs, middleware vendors, application developers
- **Many industry-leading experts involved in OpenCL's design**
  - A healthy diversity of industry perspectives
- **Apple made initial proposal and is very active in the working group**
  - Serving as specification editor

Credit: Khronos Group

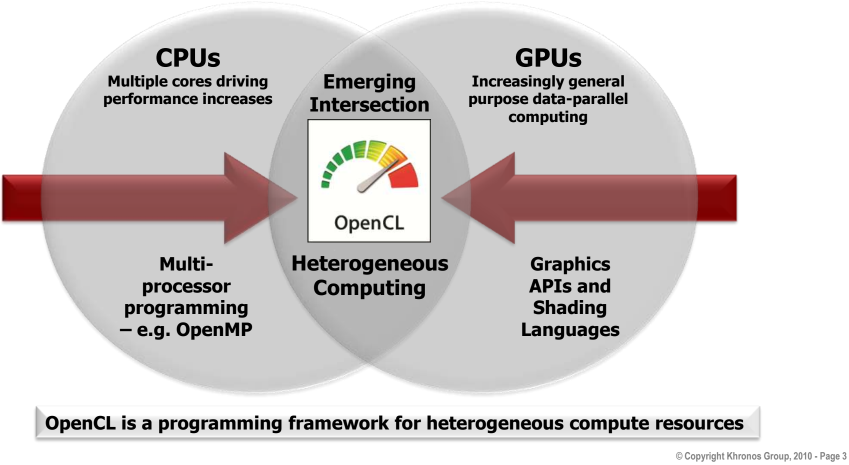
When?

- **Six months from proposal to released OpenCL 1.0 specification**
  - Due to a strong initial proposal and a shared commercial incentive
- **Multiple conformant implementations shipping**
  - Apple's Mac OS X Snow Leopard now ships with OpenCL
- **18 month cadence between OpenCL 1.0 and OpenCL 1.1**
  - Backwards compatibility protect software investment



Credit: Khronos Group

Why?



Credit: Khronos Group

CL vs CUDA side-by-side

CUDA source code:

```
__global__ void transpose(
    float *A_t, float *A,
    int a_width, int a_height)
{
  int base_idx_a     =
    blockIdx.x * BLK_SIZE +
    blockIdx.y * A_BLOCK_STRIDE;
  int base_idx_a_t   =
    blockIdx.y * BLK_SIZE +
    blockIdx.x * A_T_BLOCK_STRIDE;

  int glob_idx_a   =
    base_idx_a + threadIdx.x
    + a_width * threadIdx.y;
  int glob_idx_a_t   =
    base_idx_a_t + threadIdx.x
    + a_height * threadIdx.y;

  __shared__ float A_shared[BLK_SIZE][BLK_SIZE+1];

  A_shared[threadIdx.y][threadIdx.x] =
    A[glob_idx_a];

  __syncthreads();

  A_t[glob_idx_a_t] =
    A_shared[threadIdx.x][threadIdx.y];
}
```

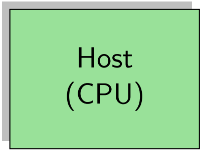
OpenCL source code:

```
void transpose(
  __global float *a_t, __global float *a,
  unsigned a_width, unsigned a_height)
{
  int base_idx_a     =
    get_group_id(0) * BLK_SIZE +
    get_group_id(1) * A_BLOCK_STRIDE;
  int base_idx_a_t   =
    get_group_id(1) * BLK_SIZE +
    get_group_id(0) * A_T_BLOCK_STRIDE;

  int glob_idx_a   =
    base_idx_a + get_local_id(0)
    + a_width * get_local_id(1);
  int glob_idx_a_t   =
    base_idx_a_t + get_local_id(0)
    + a_height * get_local_id(1);

  __local float a_local[BLK_SIZE][BLK_SIZE+1];

  a_local[get_local_id(1)*BLK_SIZE+get_local_id(0)] =
    a[glob_idx_a];

  barrier(CLK_LOCAL_MEM_FENCE);

  a_t[glob_idx_a_t] =
    a_local[get_local_id(0)*BLK_SIZE+get_local_id(1)];
}
```

OpenCL ↔ CUDA: A dictionary

OpenCL	CUDA
Grid	Grid
Work Group	Block
Work Item	Thread
`__kernel`	`__global__`
`__global`	`__device__`
`__local`	`__shared__`
`__private`	`__local__`
image*n*d_t	texture<type, *n*, ...>
`barrier(LMF)`	`__syncthreads()`
`get_local_id(012)`	`threadIdx.xyz`
`get_group_id(012)`	`blockIdx.xyz`
`get_global_id(012)`	– (reimplement)

OpenCL: Computing as a Service



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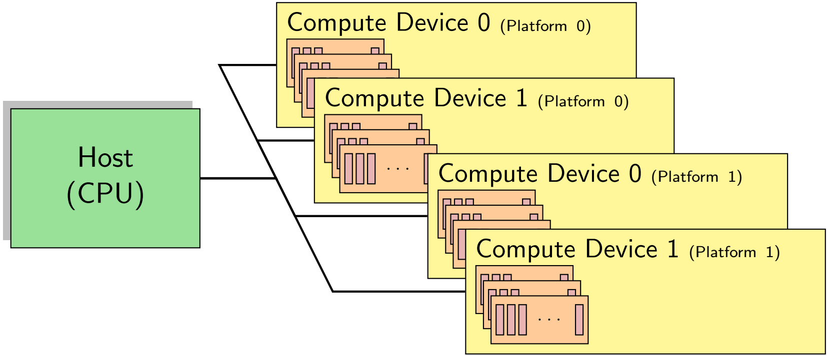

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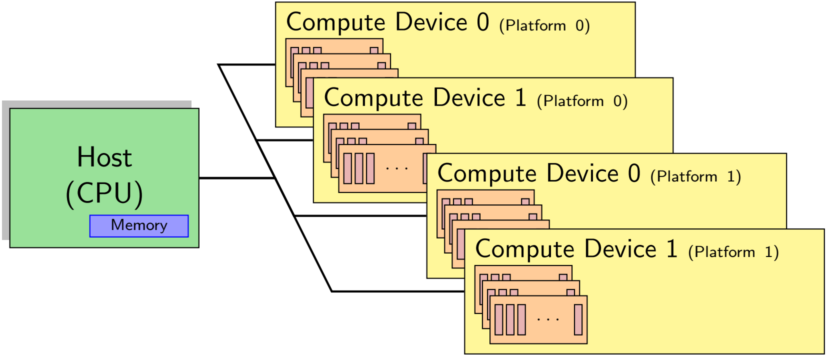

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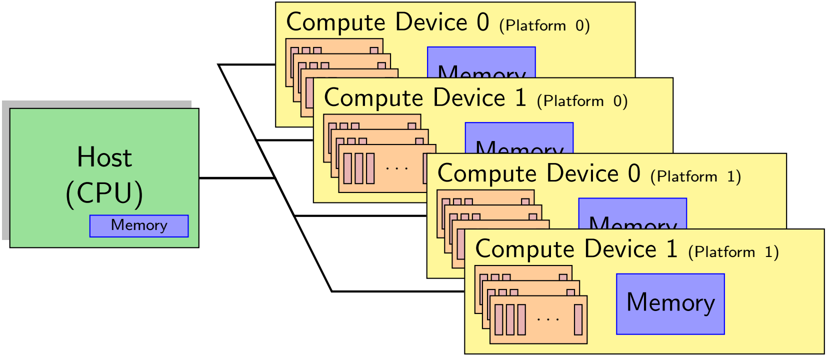

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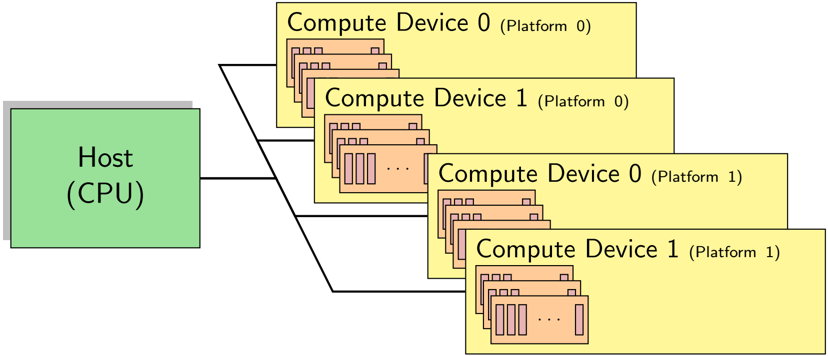

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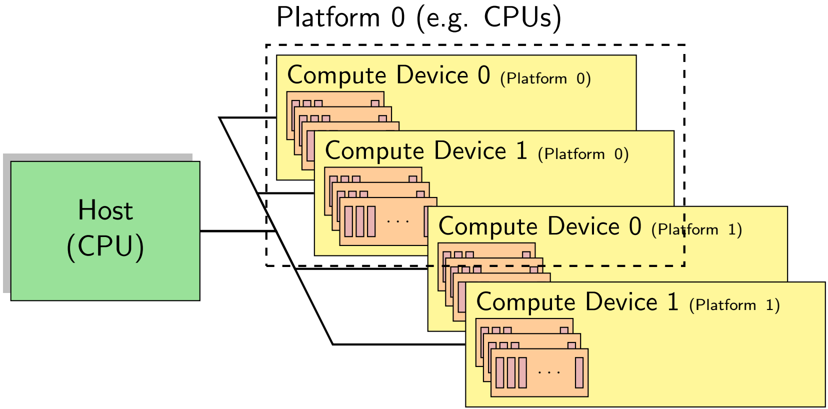

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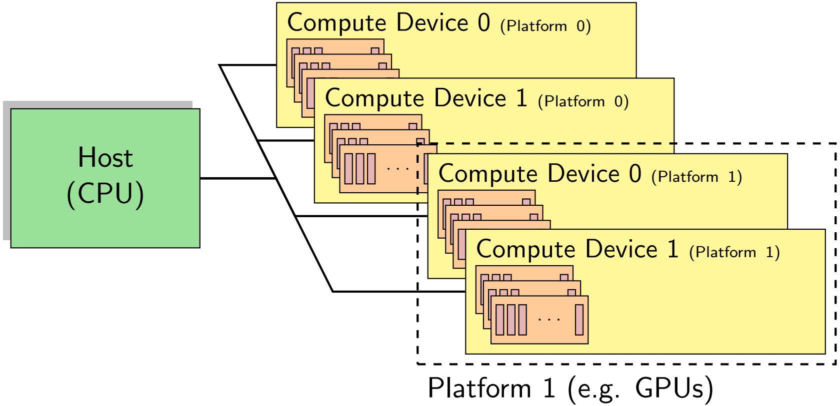

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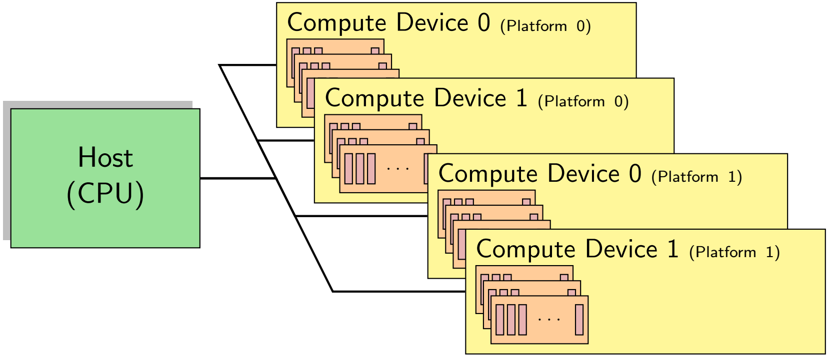

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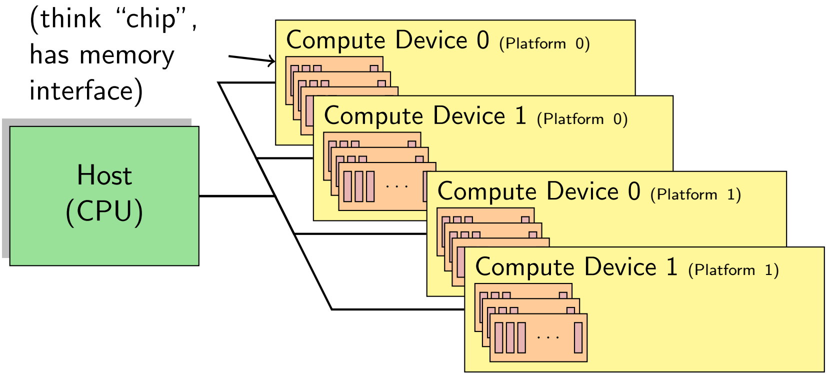

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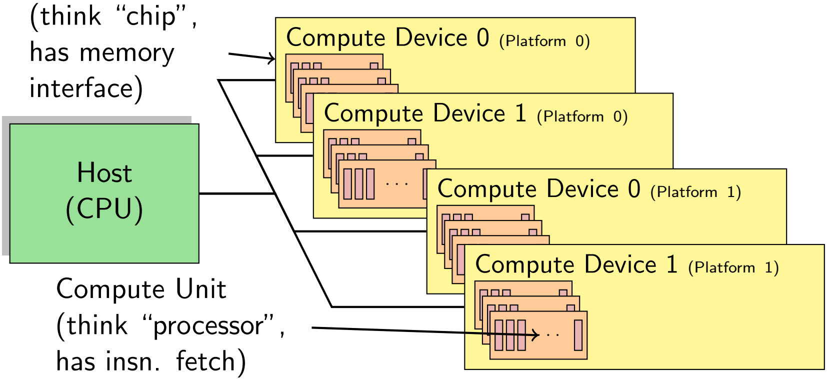

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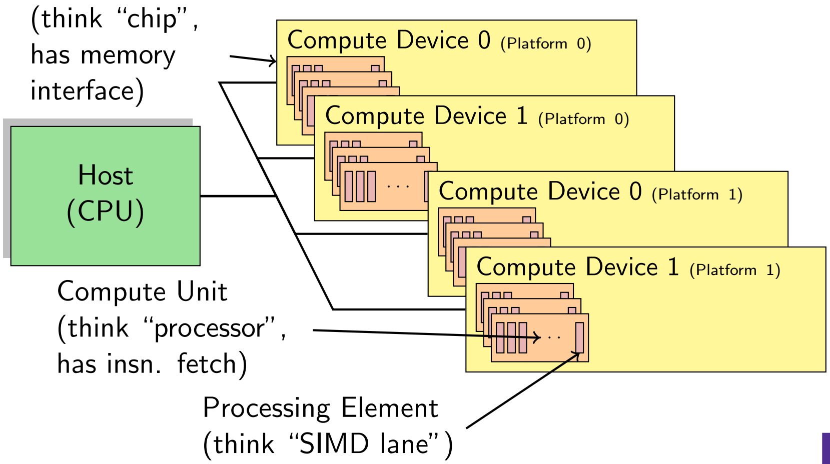

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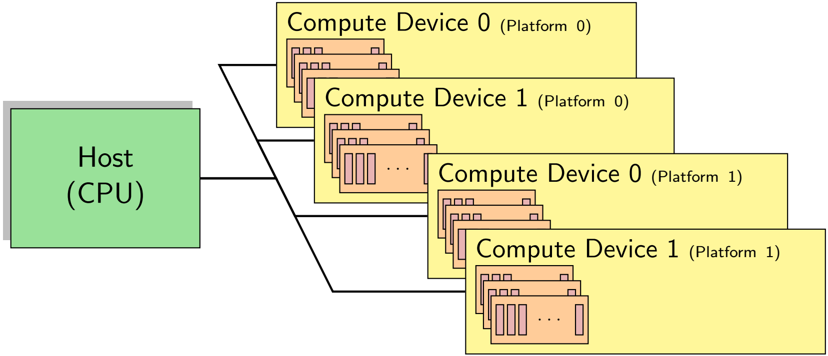

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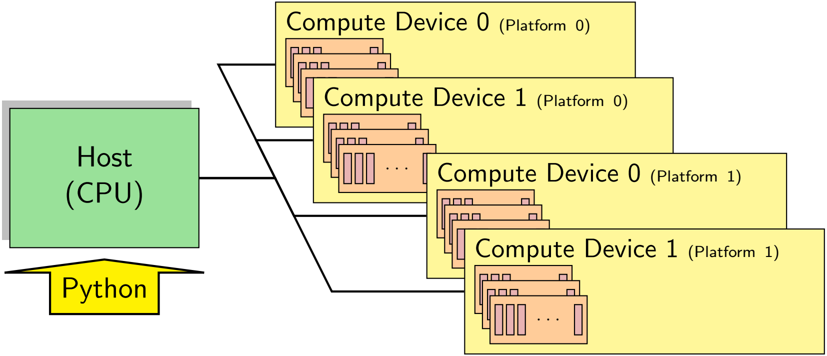

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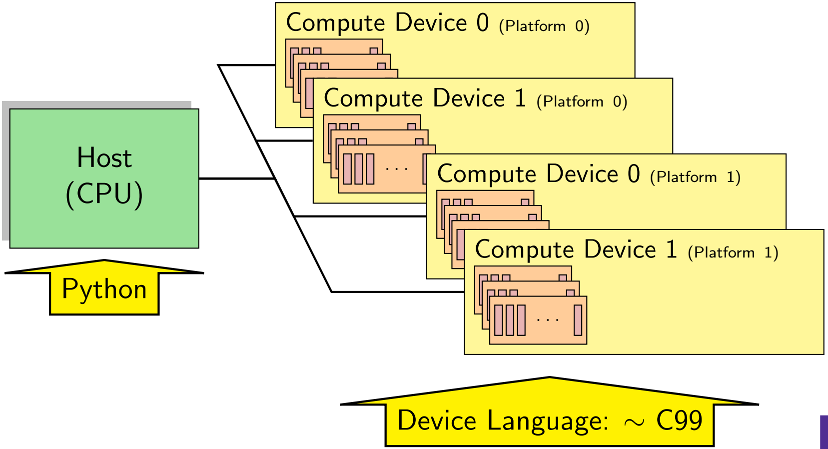

Why do Scripting for GPUs?

- GPUs are everything that scripting languages are not.
  - Highly parallel
  - Very architecture-sensitive
  - Built for maximum FP/memory throughput

  $\rightarrow$ complement each other
- CPU: largely restricted to control tasks (∼1000/sec)
  - Scripting fast enough
- Python + CUDA = **PyCUDA**
- Python + OpenCL = **PyOpenCL**

Outline

1 Intro: Python, Numpy, GPUs, OpenCL

2 GPU Programming with PyOpenCL
- First Contact
- About PyOpenCL

3 OpenCL viewed from Python

4 OpenCL implementations

Outline

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Dive into PyOpenCL

```
import pyopencl as cl, numpy

a = numpy.random.rand(256**3).astype(numpy.float32)

ctx = cl.create_some_context()
queue = cl.CommandQueue(ctx)

a_dev = cl.Buffer(ctx, cl.mem_flags.READ_WRITE, size=a.nbytes)
cl.enqueue_copy(queue, a_dev, a)

prg = cl.Program(ctx, """
    __kernel void twice(__global float *a)
    { a[get_global_id(0)] *= 2; }
    """).build()

prg.twice(queue, a.shape, (1,), a_dev)
```

Dive into PyOpenCL

```python
import pyopencl as cl, numpy

a = numpy.random.rand(256**3).astype(numpy.float32)

ctx = cl.create_some_context()
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    __kernel void twice(__global float *a)
    { a[get_global_id(0)] *= 2; }
    """).build()

prg.twice(queue, a.shape, (1,), a_dev)
```

Compute kernel

Dive into PyOpenCL: Getting Results

```python
a_dev = cl.Buffer(ctx, cl.mem_flags.READ_WRITE, size=a.nbytes)
cl.enqueue_copy(queue, a_dev, a)

prg = cl.Program(ctx, """
    __kernel void twice(__global float *a)
    { a[get_global_id(0)] *= 2; }
    """).build()

prg.twice(queue, a.shape, (1,), a_dev)

result = numpy.empty_like(a)
cl.enqueue_copy(queue, result, a_dev)
import numpy.linalg as la
assert la.norm(result - 2*a) == 0
```

Dive into PyOpenCL: Grouping

```
a_dev = cl.Buffer(ctx, cl.mem_flags.READ_WRITE, size=a.nbytes)
cl.enqueue_copy(queue, a_dev, a)

prg = cl.Program(ctx, """
    __kernel void twice(__global float *a)
    { a[get_local_id(0)+ get_local_size(0)*get_group_id(0)] *= 2; }
    """).build()

prg.twice(queue, a.shape, (256,), a_dev)

result = numpy.empty_like(a)
cl.enqueue_copy(queue, result, a_dev)
import numpy.linalg as la
assert la.norm(result - 2*a) == 0
```

Dive into PyOpenCL: Thinking on your feet

Thinking about GPU programming

How would we modify the program to...

Dive into PyOpenCL: Thinking on your feet

Thinking about GPU programming

How would we modify the program to...

1. ...compute $c_i = a_i b_i$?

Dive into PyOpenCL: Thinking on your feet

Thinking about GPU programming

How would we modify the program to. . .

1. . . . compute $c_i = a_i b_i$?
2. . . . use groups of 16×16 work items?

Dive into PyOpenCL: Thinking on your feet

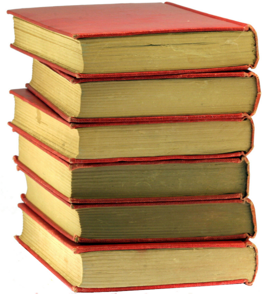
Thinking about GPU programming

How would we modify the program to...

1. ...compute $c_i = a_i b_i$?
2. ...use groups of 16×16 work items?
3. ...benchmark 1 work item per group against 256 work items per group? (Use `time.time()` and `.wait()`.)

Outline

PyOpenCL Philosophy



- Provide complete access
- Automatically manage resources
- Provide abstractions
- Allow interactive use
- Check for and report errors automatically
- Integrate tightly with `numpy`

PyOpenCL: Completeness

PyOpenCL exposes *all* of OpenCL.

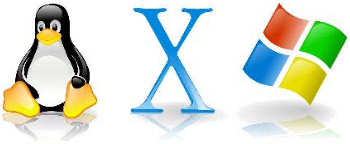
For example:

- Every `GetInfo()` query
- Images and Samplers
- Memory Maps
- Profiling and Synchronization
- GL Interop

PyOpenCL: Completeness

PyOpenCL supports (nearly) every OS that has an OpenCL implementation.

- Linux
- OS X
- Windows



Automatic Cleanup

- Reachable objects (memory, streams, ...) are never destroyed.
- Once unreachable, released at an unspecified future time.
- Scarce resources (memory) can be explicitly freed. (`obj.release()`)
- Correctly deals with multiple contexts and dependencies. (based on OpenCL's reference counting)

PyOpenCL: Documentation

PyOpenCL v0.91.2 documentation » next | modules | index

Table Of Contents

Next topic

Installation

This Page

Show Source

Quick search

Go

Enter search terms or a module, class or function name.

Welcome to PyOpenCL's documentation!

PyOpenCL gives you easy, Pythonic access to the OpenCL parallel computation API. What makes PyOpenCL special?

- Object cleanup tied to lifetime of objects. This idiom, often called RAII in C++, makes it much easier to write correct, leak- and crash-free code.
- Completeness. PyOpenCL puts the full power of OpenCL's API at your disposal, if you wish. Every obscure *get_info()* query and all CL calls are accessible.
- Automatic Error Checking. All errors are automatically translated into Python exceptions.
- Speed. PyOpenCL's base layer is written in C++, so all the niceties above are virtually free.
- Helpful Documentation. You're looking at it. ;)
- Liberal license. PyOpenCL is open-source under the *MIT license* and free for commercial, academic, and private use.

Here's an example, to give you an impression:

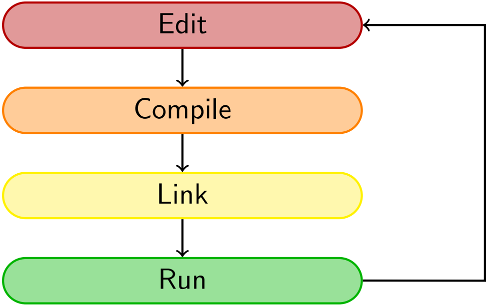
```
import pyopencl as cl
import numpy
import numpy.linalg as la

a = numpy.random.rand(50000).astype(numpy.float32)
b = numpy.random.rand(50000).astype(numpy.float32)

ctx = cl.Context()
queue = cl.CommandQueue(ctx)

mf = cl.mem_flags
a_buf = cl.Buffer(ctx, mf.READ_ONLY | mf.COPY_HOST_PTR, hostbuf=a)
b_buf = cl.Buffer(ctx, mf.READ_ONLY | mf.COPY_HOST_PTR, hostbuf=b)
dest_buf = cl.Buffer(ctx, mf.WRITE_ONLY, b.nbytes)

prg = cl.Program(ctx, """
    __kernel void sum(__global const float *a,
    __global const float *b, __global float *c)
    {
      int gid = get_global_id(0);
      c[gid] = a[gid] + b[gid];
    }
    """).build()

prg.sum(queue, a.shape, a_buf, b_buf, dest_buf)

a_plus_b = numpy.empty_like(a)
cl.enqueue_read_buffer(queue, dest_buf, a_plus_b).wait()

print la.norm(a_plus_b - (a+b))
```

(You can find this example as `examples/demo.py` in the PyOpenCL source distribution.)

Contents

Scripting: Interpreted, not Compiled

Program creation workflow:



Scripting: Interpreted, not Compiled

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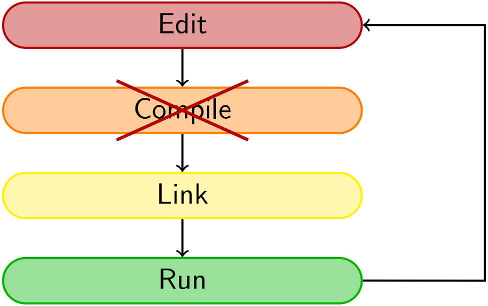

Scripting: Interpreted, not Compiled

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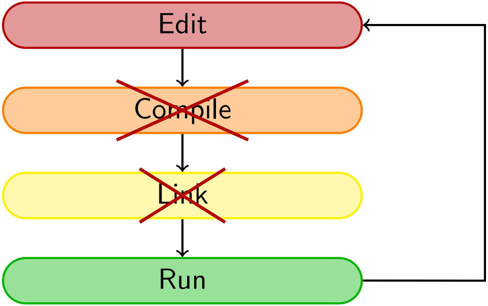
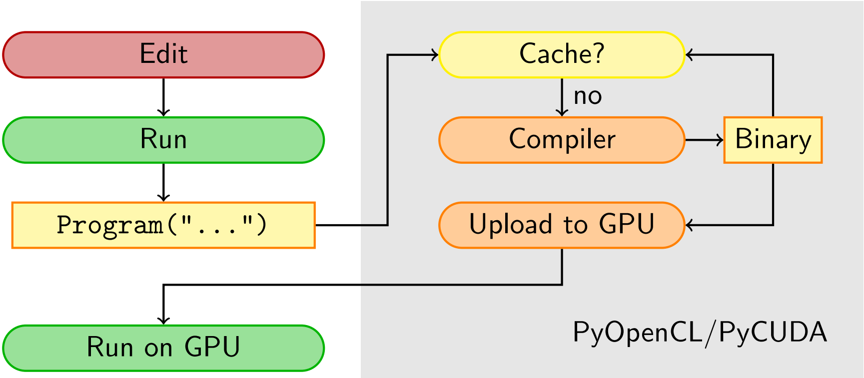


PyOpenCL, PyCUDA: Workflow



PyOpenCL: Vital Information

- `http://mathema.tician.de/software/pyopencl`
  - Downloaded 30k+ times
- Complete documentation
- MIT License (no warranty, free for all use)
- Requires: numpy, Python 2.4+.
- Community: mailing list, wiki
- Add-on packages (e.g. PyFFT, Sailfish, PyWENO)

An Appetizer

Remember your first PyOpenCL program?

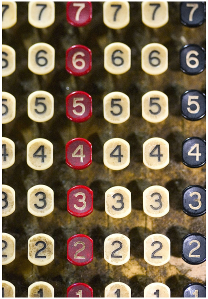
Abstraction is good:

```
import numpy
import pyopencl as cl
import pyopencl.array as cl_array

ctx = cl.create_some_context()
queue = cl.CommandQueue(ctx)

a_gpu = cl_array.to_device(
        ctx, queue, numpy.random.randn(4,4).astype(numpy.float32))
a_doubled = (2*a_gpu).get()
print a_doubled
print a_gpu
```

pyopencl.array: Simple Linear Algebra

pyopencl.array.Array:

- Meant to look and feel just like numpy.
  - p.a.to_device(ctx, queue, numpy_array)
  - numpy_array = ary.get()
- +, -, ∗, /, fill, sin, arange, exp, rand, ...
- Mixed types (int32 + float32 = float64)
- print cl_array for debugging.
- Allows access to raw bits
  - Use as kernel arguments, memory maps



pyopencl.elementwise: Elementwise expressions

Avoiding extra store-fetch cycles for elementwise math:

```
n = 10000
a_gpu = cl_array.to_device(
        ctx, queue, numpy.random.randn(n).astype(numpy.float32))
b_gpu = cl_array.to_device(
        ctx, queue, numpy.random.randn(n).astype(numpy.float32))

from pyopencl.elementwise import ElementwiseKernel
lin_comb = ElementwiseKernel(ctx,
        "float a, float *x, float b, float *y, float *z",
        "z[i] = a*x[i] + b*y[i]")

c_gpu = cl_array.empty_like(a_gpu)
lin_comb(5, a_gpu, 6, b_gpu, c_gpu)

import numpy.linalg as la
assert la.norm((c_gpu - (5*a_gpu+6*b_gpu)).get()) < 1e-5
```

pyopencl.reduction: Reduction made easy

Example: A dot product calculation

```
from pyopencl.reduction import ReductionKernel
dot = ReductionKernel(ctx, dtype_out=numpy.float32, neutral="0",
    reduce_expr="a+b", map_expr="x[i]*y[i]",
    arguments="__global const float *x, __global const float *y")

import pyopencl.clrandom as cl_rand
x = cl_rand.rand(ctx, queue, (1000*1000), dtype=numpy.float32)
y = cl_rand.rand(ctx, queue, (1000*1000), dtype=numpy.float32)

x_dot_y = dot(x, y).get()
x_dot_y_cpu = numpy.dot(x.get(), y.get())
```

pyopencl.scan: Scan made easy

Example: A cumulative sum computation

```
from pyopencl.scan import InclusiveScanKernel
knl = InclusiveScanKernel(ctx, np.int32, "a+b")

n = 2**20-2**18+5
host_data = np.random.randint(0, 10, n).astype(np.int32)
dev_data = cl_array.to_device(queue, host_data)

knl(dev_data)
assert (dev_data.get() == np.cumsum(host_data, axis=0)).all()
```

Questions?

?

Outline

1 Intro: Python, Numpy, GPUs, OpenCL

2 GPU Programming with PyOpenCL

3 OpenCL viewed from Python
- Device Language
- The OpenCL runtime
- Synchronization
- Extensions

4 OpenCL implementations

Measuring Performance

Writing high-performance Codes

Mindset: What is going to be the limiting factor?

- Floating point throughput?
- Memory bandwidth?
  - Cache sizes?

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Benchmark the assumed limiting factor right away.

Measuring Performance

Writing high-performance Codes

Mindset: What is going to be the limiting factor?

- Floating point throughput?
- Memory bandwidth?
    - Cache sizes?

Benchmark the assumed limiting factor right away.

Evaluate

- Know your peak throughputs (roughly)
- Are you getting close?
- Are you tracking the right limiting factor?

Outline

OpenCL Device Language

OpenCL device language is C99, with these differences:

- Index getters
- Memory space qualifiers
- Vector data types
- Many generic ('overloaded') math functions including fast `native_...` varieties.
- Synchronization
- Recursion
- `malloc()`

Address Space Qualifiers

Type	Per	Access	Latency	
private	work item	R/W	1 or 1000	
local	group	R/W	2	
global	grid	R/W	1000	Cached?
constant	grid	R/O	1-1000	Cached
image*n*d_t	grid	R(/W)	1000	Spatially cached

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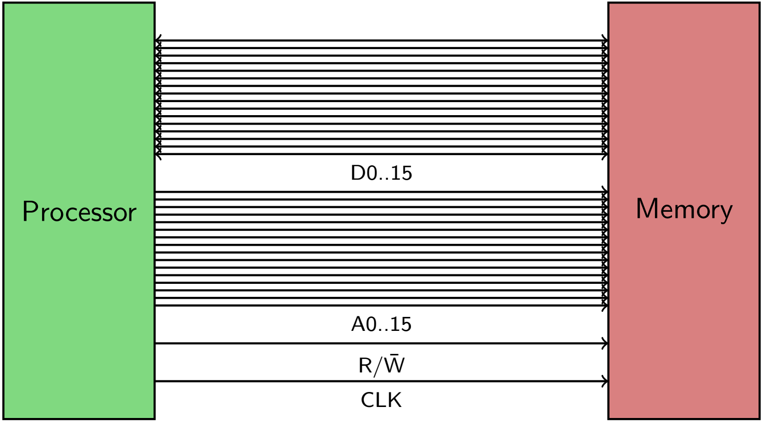
Important

Different types of memory are good at different types of access. Successful algorithms combine many types' strengths.

How does computer memory work?

One (reading) memory transaction (simplified):



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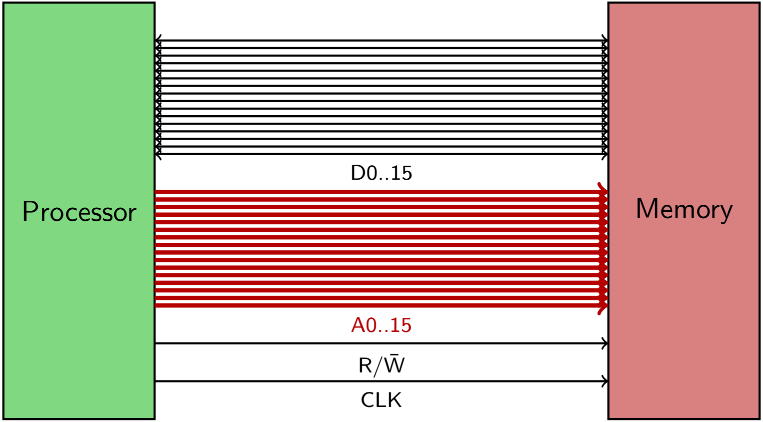

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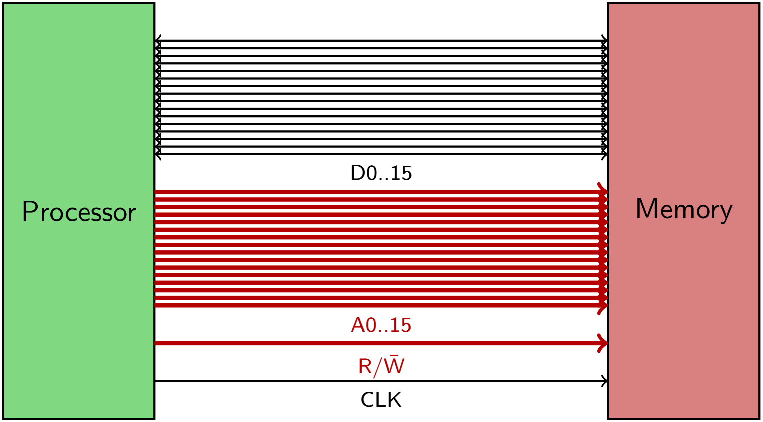

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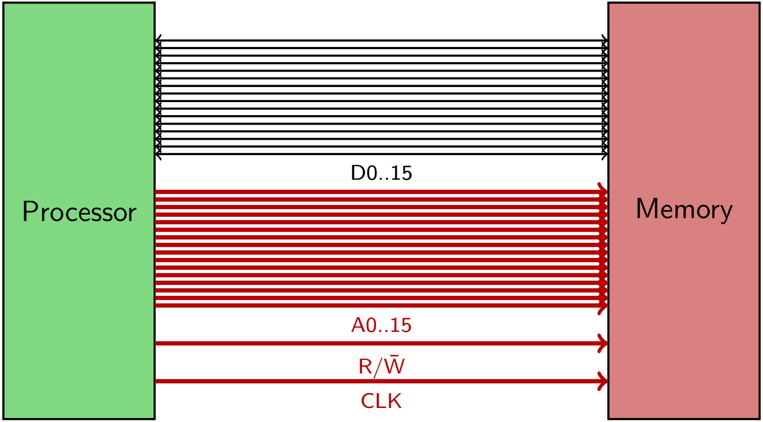

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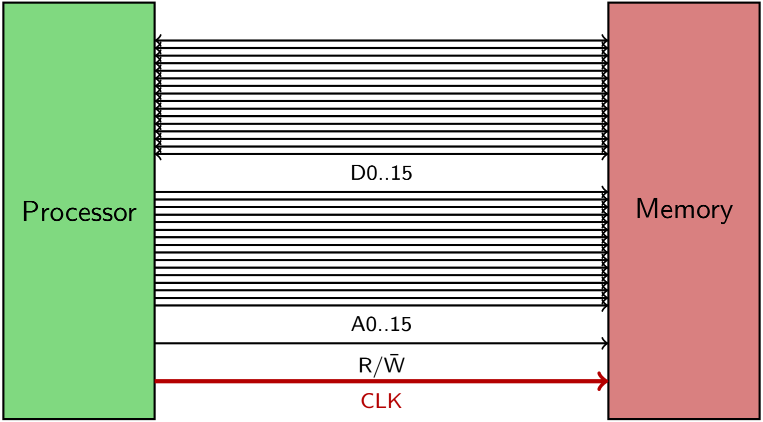

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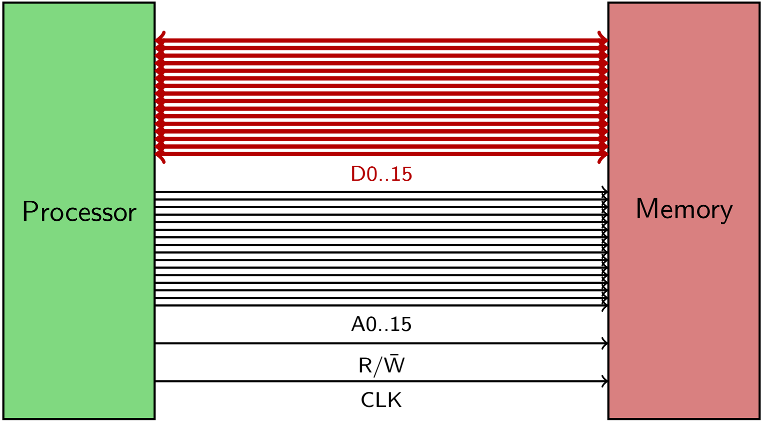

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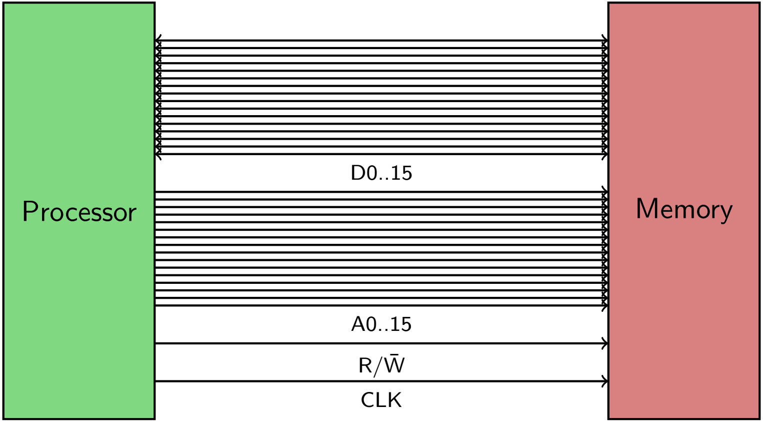


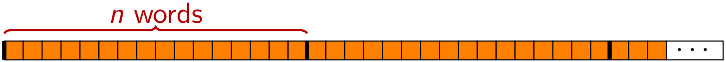
Observation: Access (and addressing) happens in bus-width-size "chunks".

Global Memory

Rule of thumb

$$n = \min\left(\frac{\text{Bus width in bits}}{\text{Word size in bits}}, \text{SIMD group size}\right)$$

work items access global memory simultaneously. Full utilization only if all bits in bus transaction are useful.



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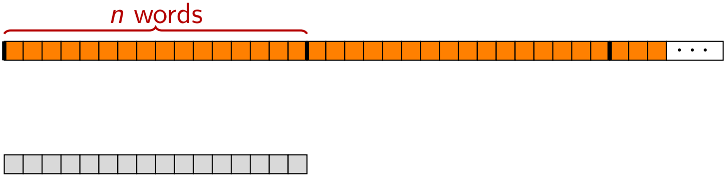

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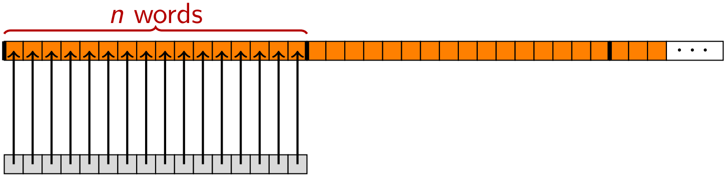

OK: `global_variable[get_global_id(0)]`
(Single transaction)

Global Memory

Rule of thumb

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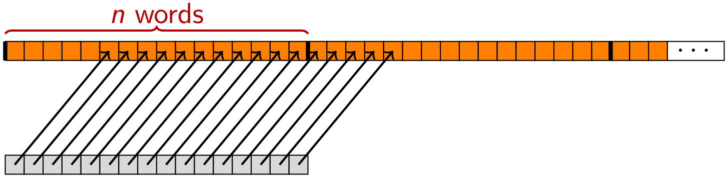

Bad: `global_variable[5+get_global_id(0)]`
(Two transactions)

Global Memory

Rule of thumb

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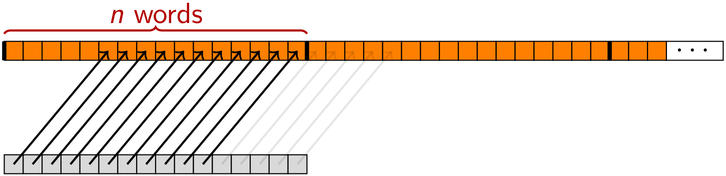

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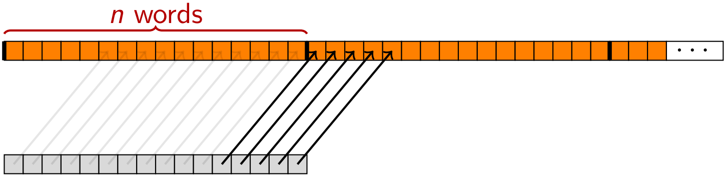

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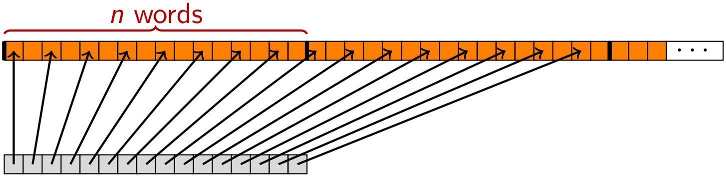


Bad: `global_variable[2*get_global_id(0)]`
(Two transactions)

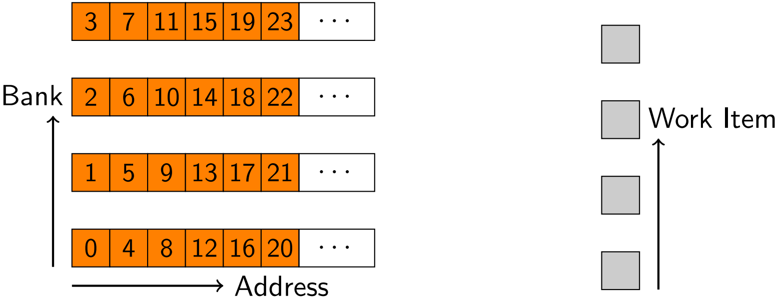
Making sense of Global Memory

Consider the following examples:

- List of XYZ vectors:
  - XXXX...YYYY...ZZZZ...(“SoA”)
  - XYZXYZXYZ...(“AoS”)
- Accessing a row-major (C order) matrix
  - by rows
  - by columns

Local Memory: Banking

Local Memory: Banking



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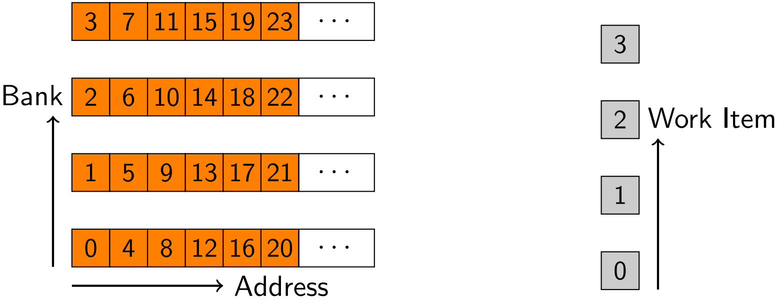
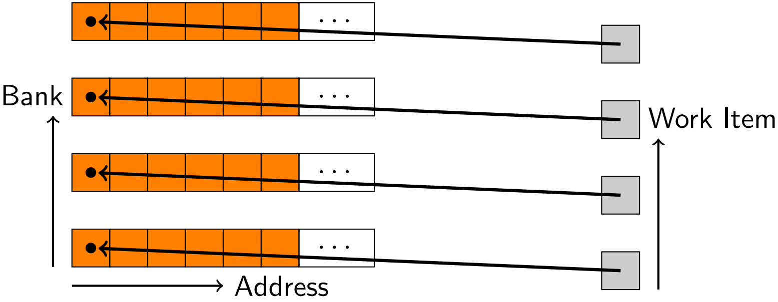
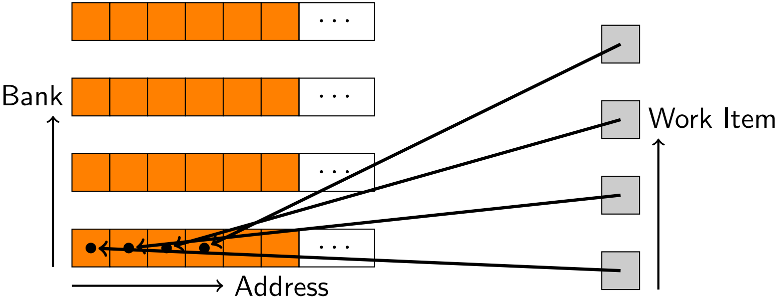


Local Memory: Banking



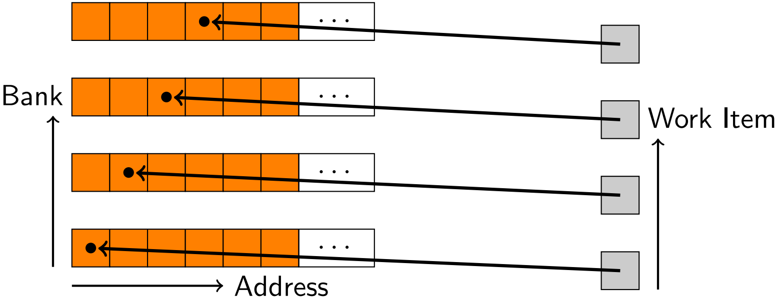
OK: `local_variable[get_local_id(0)]`,
(Single cycle)

Local Memory: Banking



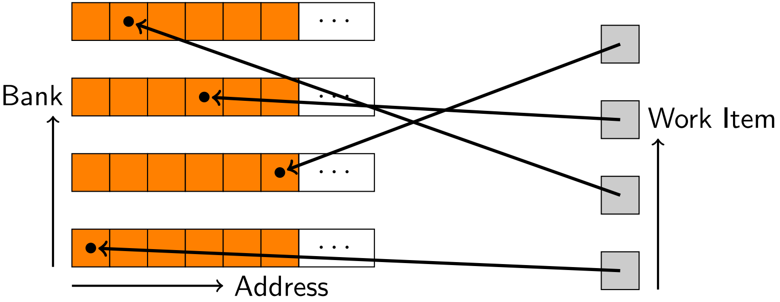
Bad: `local_variable[BANK_COUNT*get_local_id(0)]`
(`BANK_COUNT` cycles)

Local Memory: Banking



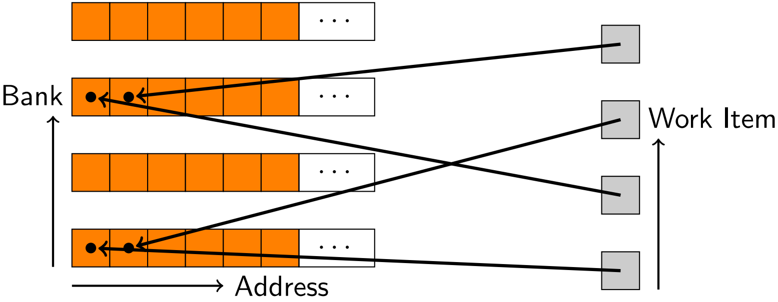
OK: `local_variable[(BANK_COUNT+1)*get_local_id(0)]`
(Single cycle)

Local Memory: Banking



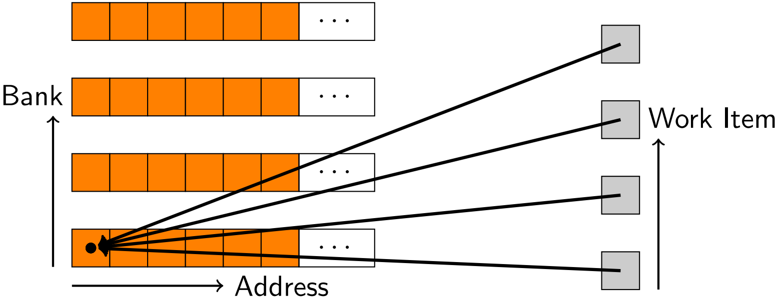
OK: `local_variable[ODD_NUMBER*get_local_id(0)]`
(Single cycle)

Local Memory: Banking



Bad: `local_variable[2*get_local_id(0)]`
(`BANK_COUNT`/2 cycles)

Local Memory: Banking



OK: `local_variable[f(get_group_id(0))]`
(Broadcast–single cycle)

Local Memory: Banking

Example: Nvidia GT200 has 16 banks.
Work items access local memory in groups of 16.

CL vector data types

`float`*n* `vec` (*n*=1,2,3,4,8,16) (also for `double` and integer types) Components:

- `vec.s012...abcdef` (or `xyzw`)
- `vec.s3120` (Swizzling)
- `vec.s024 = (float3)(1,2,3);` (Lvalue, Literals)

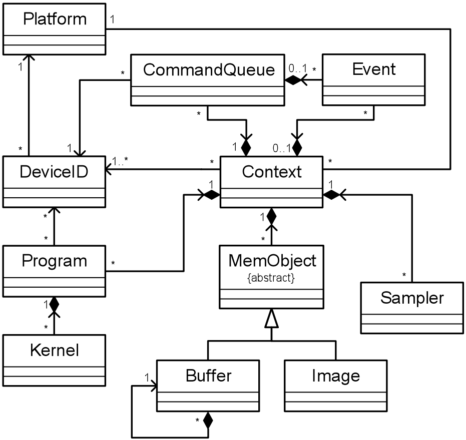
Usage:

- Elementwise operations (+,-,`sin` (generic!),...)
- `float`*n* `vload`*n*`/vstore`*n*`(offset, float *)` (aligned!)
- `dot/distance`

Using CPU implementation: One of the sanest ways of using SSE/vector intrinsics!

Outline

OpenCL Object Diagram



Credit: Khronos Group

CL "Platform"

- "Platform": a collection of devices, all from the same *vendor*.
- All devices in a platform use same CL driver/implementation.
- Multiple platforms can be used from one program → *ICD*.

`libOpenCL.so`: ICD loader

`/etc/OpenCL/vendors/`*`somename`*`.icd`: Plain text file with name of `.so` containing CL implementation.

CL "Compute Device"

CL Compute Devices:

- CPUs, GPUs, accelerators, ...
  - Anything that fits the programming model.
- A processor die with an interface to off-chip memory
- Can get list of devices from platform.

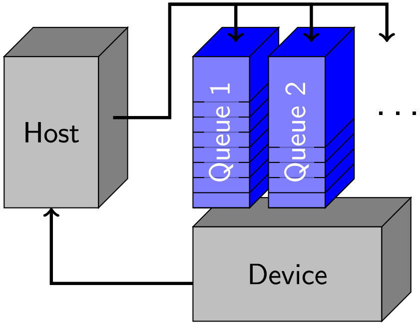
Contexts

```
context = cl.Context(devices=None | [dev1, dev2], dev_type=None)
context = cl.create_some_context( interactive =True)
```

- Spans one or more Devices
- Create from device type or list of devices
  - See docs for `cl.Platform`, `cl.Device`
- `dev_type`: *`DEFAULT`*, `ALL`, `CPU`, `GPU`
- Needed to...
  - ...allocate Memory Objects
  - ...create and build Programs
  - ...host Command Queues
  - ...execute Grids

OpenCL: Command Queues

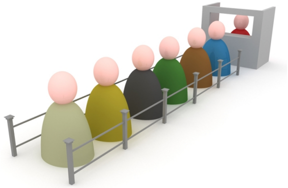
- Host and Device run asynchronously
- Host submits to queue:
  - Computations
  - Memory Transfers
  - Sync primitives
  - ...
- Host can wait for drained queue
- Profiling



Command Queues and Events

```
queue = cl.CommandQueue(context, device=None,
    properties=None | [(prop, value ),...])
```

- Attached to single device
- cl.command_queue_properties...
  - `OUT_OF_ORDER_EXEC_MODE_ENABLE`: Do not force sequential execution
  - `PROFILING_ENABLE`: Gather timing info



Building Blocks in Action

```
import pyopencl as cl

platforms = cl.get_platforms()
my_platform = platforms[0]
print my_platform.vendor

devices = my_platform.get_devices()
my_device = devices[0]
print my_device.name

ctx = cl.Context([my_device])

cpq = cl.command_queue_properties
queue = cl.CommandQueue(ctx, my_device, cpq.PROFILING_ENABLE)
```

Simple version:

```
ctx2 = cl.create_some_context()
queue2 = cl.CommandQueue(ctx)
```

Command Queues and Events

```
event = cl.enqueue_XXX(queue, ..., wait_for=[evt1, evt2])
```

Every enqueue operation returns an *Event*.

Also possible: Operation-less events ("Markers")

- Wait (`evt.wait()`), polling
- Specify dependencies

  Every enqueue operation takes a list arg `wait_for` of dependencies.
- Profile

  `event.profile....`
  - `QUEUED`, `SUBMIT`
  - `START`, `END`

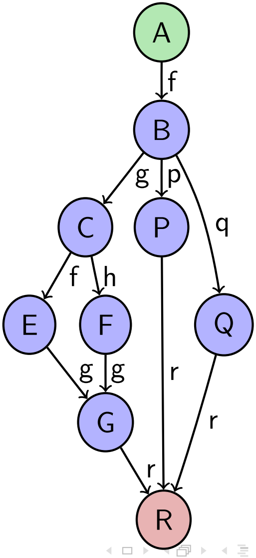
  (time stamp in ns)

Profiling example

```
start_event = cl.enqueue_marker(queue)

# enqueue some commands

stop_event = cl.enqueue_marker(queue)
stop_event.wait()

elapsed_seconds = 1e-9*(
        start_event.profile.END - start_event.profile.END)

# --- OR ---

op_event = knl(queue, global_size, grp_size, args ...)
op_event.wait()
elapsed_seconds = 1e-9*(
        op_event.profile.END - op_event.profile.START)
```

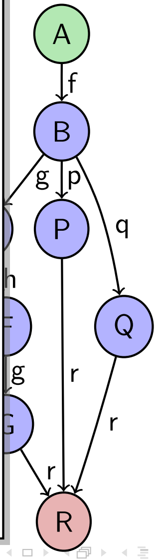
Capturing Dependencies

$B = f(A)$
$C = g(B)$
$E = f(C)$
$F = h(C)$
$G = g(E,F)$
$P = p(B)$
$Q = q(B)$
$R = r(G,P,Q)$



Ca

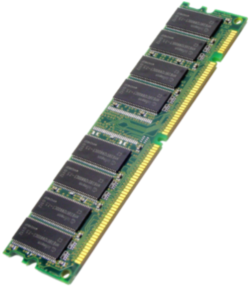
- Switch queue to out-of-order mode!
- Specify as list of events using `wait_for=` optional keyword to `enqueue_XXX`.
- Can also enqueue barrier.
- Common use case: Transmit/receive from other MPI ranks.
- Possible in hardware on Nv Fermi, AMD Cayman: Submit parallel work to increase machine use.
  - Not yet ubiquitously implemented



Memory Objects: Buffers

```
buf = cl.Buffer(context, flags, size=0, hostbuf=None)
```

- Chunk of device memory
- No type information: “Bag of bytes”
- Observe: *Not* tied to device.
  → no fixed memory address
  → pointers do *not* survive kernel launches
  → movable between devices
- `flags`:
  - `READ_ONLY/WRITE_ONLY/READ_WRITE`
  - `{ALLOC,COPY,USE}_HOST_PTR`



Memory Objects: Buffers

```
buf = cl.Buffer(context, flags, size=0, hostbuf=None)
```

`COPY_HOST_PTR`:

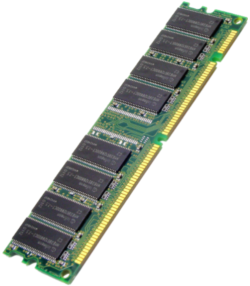
- Use `hostbuf` as initial content of buffer

`USE_HOST_PTR`:

- `hostbuf` *is* the buffer.
- Caching in device memory is allowed.

`ALLOC_HOST_PTR`:

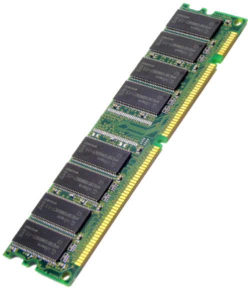
- *New* host memory (unrelated to `hostbuf`) is visible from device *and* host.



Memory Objects: Buffers

```
buf = cl.Buffer(context, flags, size=0, hostbuf=None)
```

- Specify `hostbuf` or `size` (or both)
- `hostbuf`: Needs Python Buffer Interface e.g. `numpy.ndarray`, `str`.
  - Important: Memory layout matters
- Passed to device code as pointers (e.g. `float *`, `int *`)
- enqueue_copy(queue, dest, src)
- Can be mapped into host address space: `cl.MemoryMap`.



NYU

Command Queues and Buffers: A Crashy Puzzle

✔ OK

```
a = numpy.random.rand(256**3).astype(numpy.float32)
a_dev = cl.Buffer(ctx, cl.mem_flags.READ_WRITE, size=a.nbytes)
cl.enqueue_copy(queue, a_dev, a,
        is_blocking=False)
```

Command Queues and Buffers: A Crashy Puzzle

✔ OK

```
a = numpy.random.rand(256**3).astype(numpy.float32)
a_dev = cl.Buffer(ctx, cl.mem_flags.READ_WRITE, size=a.nbytes)
cl.enqueue_copy(queue, a_dev, a,
        is_blocking=False)
```

✖ Crash

```
a_dev = cl.Buffer(ctx, cl.mem_flags.READ_WRITE, size=256**3*4)
cl.enqueue_copy(queue, a_dev,
        numpy.random.rand(256**3).astype(numpy.float32),
        is_blocking=False)
```

Command Queues and Buffers: A Crashy Puzzle

✔ OK

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a = numpy.random.rand(256**3).astype(numpy.float32)
a_dev = cl.Buffer(ctx, cl.mem_flags.READ_WRITE, size=a.nbytes)
cl.enqueue_copy(queue, a_dev, a,
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```

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        numpy.random.rand(256**3).astype(numpy.float32),
        is_blocking=False)
```

✔ OK

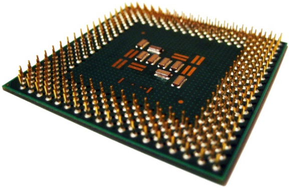
```
a_dev = cl.Buffer(ctx, cl.mem
cl.enqueue_copy(queue, a_dev,
        numpy.random.rand(25
        is_blocking=True)
```

Improved in PyOpenCL 2011.2: "nanny" events.

Programs and Kernels

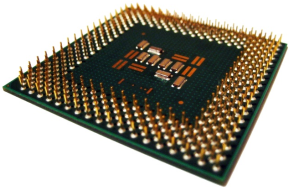
```python
prg = cl.Program(context, src)
```

- `src`: OpenCL device code
  - Derivative of C99
  - Functions with `__kernel` attribute can be invoked from host
- `prg.build(options="", devices=None)`
- `kernel = prg.kernel_name`
- `kernel(queue,` (G_x, G_y, G_z), (L_x, L_y, L_z), `arg, ..., wait_for=None)`



Program Objects

kernel (queue, (Gx,Gy,Gz), (Sx,Sy,Sz), arg, ..., wait_for =None)

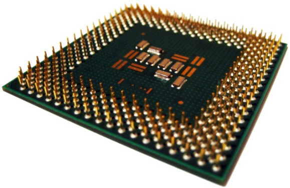


arg may be:

- `None` (a `NULL` pointer)
- numpy sized scalars: `numpy.int64,numpy.float32,...`
- Anything with buffer interface: `numpy.ndarray`, `str`
- Buffer Objects
- Also: `cl.Image`, `cl.Sampler`, `cl.LocalMemory`

Program Objects

```
kernel(queue, (Gx,Gy,Gz), (Sx,Sy,Sz), arg, ..., wait_for=None)
```



Explicitly sized scalars:

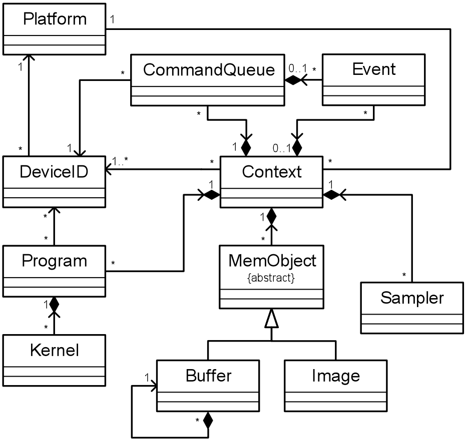
✖ Annoying, error-prone.

Better:

```
kernel.set_scalar_arg_dtypes([
   numpy.int32, None,
   numpy.float32])
```

Use `None` for non-scalars.

OpenCL Object Diagram



Credit: Khronos Group

Outline

1 Intro: Python, Numpy, GPUs, OpenCL

2 GPU Programming with PyOpenCL

3 OpenCL viewed from Python
- Device Language
- The OpenCL runtime
- Synchronization
- Extensions

4 OpenCL implementations

Recap: Concurrency and Synchronization

GPUs have layers of concurrency.

Each layer has its synchronization primitives.

Recap: Concurrency and Synchronization

GPUs have layers of concurrency.
Each layer has its synchronization primitives.

- Intra-group:
  `barrier(...)`,
  `mem_fence(...)`
  `...` =
  `CLK_{LOCAL,GLOBAL}_MEM_FENCE`
- Inter-group:
  Kernel launch
- CPU-GPU:
  Command queues, Events

Synchronization between Groups

Golden Rule:

Results of the algorithm must be independent of the order in which work groups are executed.

Synchronization between Groups

Golden Rule:

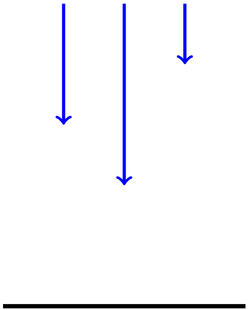
Results of the algorithm must be independent of the order in which work groups are executed.

Consequences:

- Work groups may read the same information from global memory.
- But: Two work groups may not validly write different things to the same global memory.
- Kernel launch serves as
  - Global barrier
  - Global memory fence

Synchronization

What is a Barrier?



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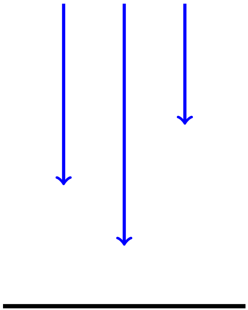

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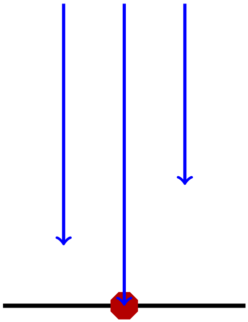

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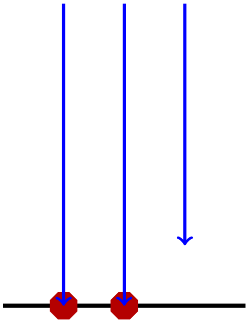

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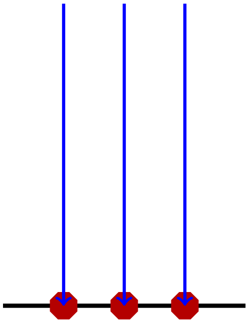

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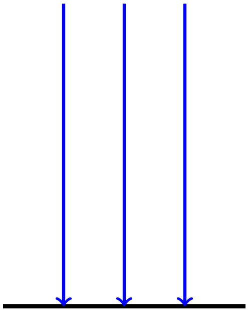

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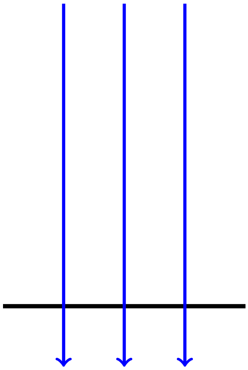

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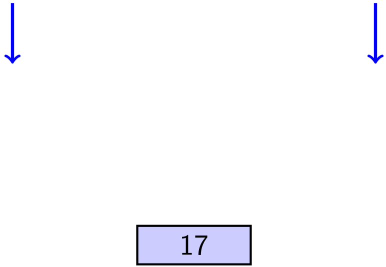
What is a Memory Fence?



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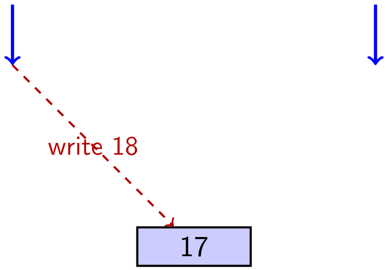

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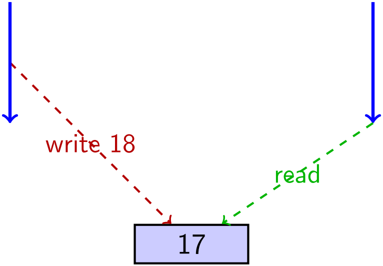

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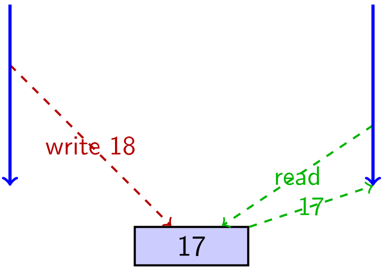

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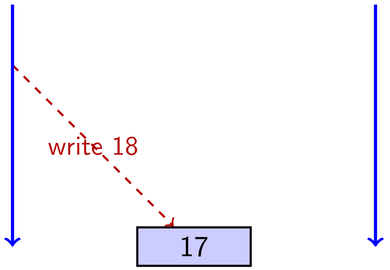

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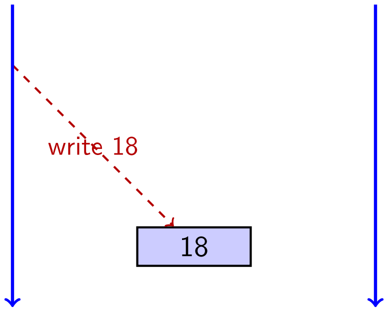

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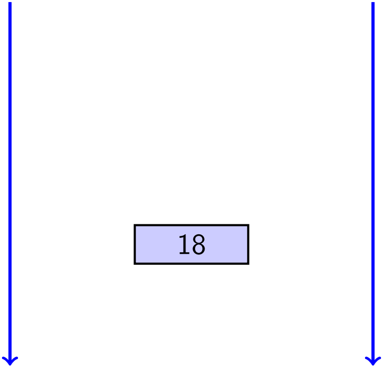

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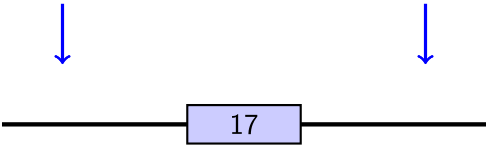

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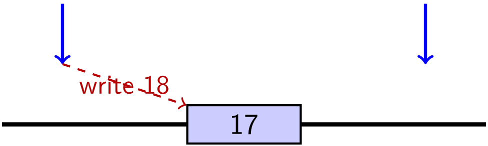

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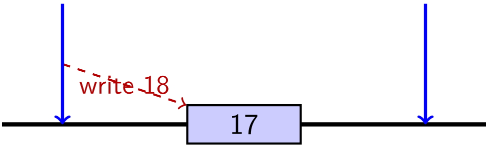

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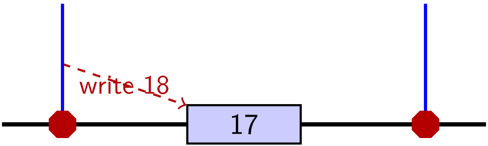

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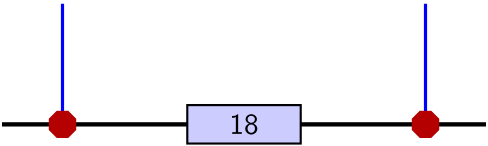

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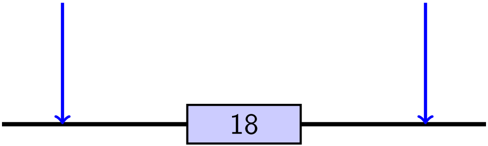

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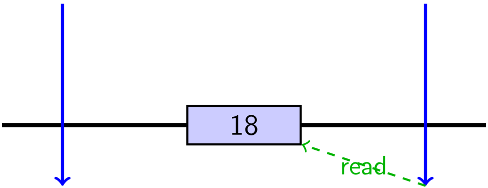

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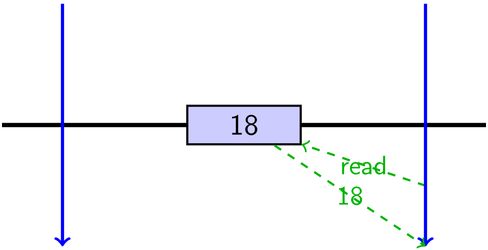

Atomic Operations

Collaborative (inter-block) Global Memory Update:

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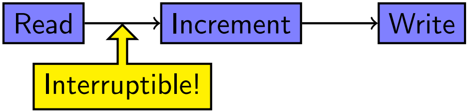

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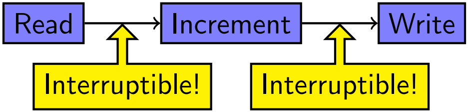

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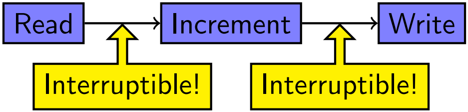

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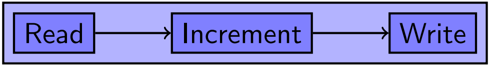

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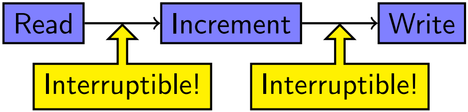

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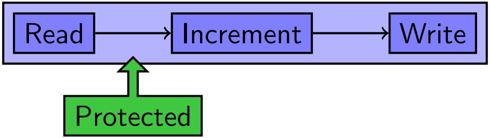

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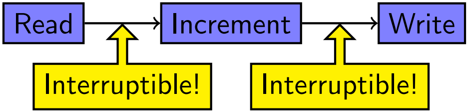

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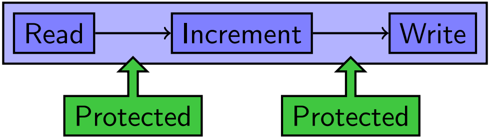

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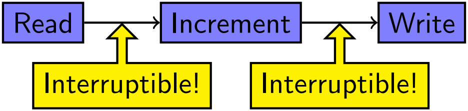

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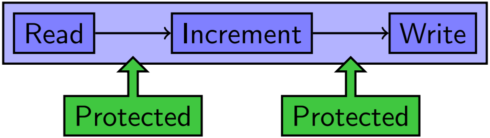


How?

atomic_{add,inc,cmpxchg,...}(int *global, int value);

Outline

Extensions: Future-proof CL

Similar extensions mechanism to OpenGL.

Two mechanisms:

- Runtime:
  - `cl_ext.h` header
  - availability checkable via `#ifdef`
  - `device.extensions`
- Device language:
  `#pragma OPENCL EXTENSION name : enable`

Important extension:

- cl_khr_fp64

Vendor and 'official' extensions.

Extension Example: `cl_ext_migrate_memobject`

- CL Memory Objects (Buffers, Images) tied to *context*, not *device*
- CL Standard: Implicit migration of data to location of use
- Compliant implementations are allowed to store all data on host, transfer out just for kernel
- With migration extension:
  - Migration becomes schedulable, takes part in command queue
  - More control over data locality
- ✔ Supported by PyOpenCL

Extension Example: cl_ext_device_fission

- Can partition a compute device
  - Equally
  - By name, counts
  - By affinity domain (L*n* cache, NUMA
- Help avoid starvation of processes that need a certain minimum throughput.
- Makes two-kernel producer-consumer model feasible.
  - Otherwise: No guarantee of progress!
- Available on Intel, AMD (CPU+GPU!)
- ✔ Supported by PyOpenCL

Outline

The Nvidia CL implementation

Targets only GPUs

Notes:

- Nearly identical to CUDA
  - No native C-level JIT in CUDA (→ PyCUDA)
- Page-locked memory:
  Use `CL_MEM_ALLOC_HOST_PTR`.
  (Careful: double meaning)
- No linear memory texturing
- CUDA device emulation mode deprecated
  → Use AMD CPU CL (faster, too!)

The Apple CL implementation

Targets CPUs and GPUs

General notes:

- Different header name `OpenCL/cl.h` instead of `CL/cl.h` Use `-framework OpenCL` for C access.
- Beware of imperfect compiler cache implementation (ignores include files)

CPU notes:

- One work item per processor

GPU similar to hardware vendor implementation.
(New: Intel w/ Sandy Bridge)

The AMD CL implementation

Targets CPUs and GPUs (from both AMD and Nvidia)

GPU notes:

- Wide SIMD groups (64)
- VLIW4 (previously VLIW5)
  - *very* flop-heavy machine
  - $\rightarrow$ ILP and explicit SIMD
  - Non-vector memory coalescing only on Cayman+
- GCN: Vector *and* scalar unit
  - Move towards Nv-like programming model

CPU notes:

- Many work items per processor (emulated)
- `cl_amd_printf`
- "APU": CPU/GPU integration not very tight yet

The Intel CL implementation

CPUs now, GPUs with Ivy Bridge+

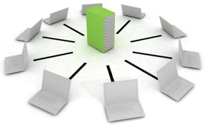
CPU notes:

- Good vectorizing compiler
- Only implementation of out-of-order queues for now
- Based on Intel TBB

GPU notes:

- Flexible design: SIMDm VLIWn
- Lots of fixed-function hardware
- Last-level Cache (LLC) integrated between CPU and GPU

The MOSIX Virtual CL implementation



- Aggregates all CL devices on a cluster into a single platform
- Looks like a "regular" CL implementation to the user
- Obvious scaling limits, but useful if the application is right
- Just heard from author: PyOpenCL supported as of version 1.10
- Aggregates communication to avoid network round-trips

Questions?

?

Image Credits

- Isaiah die shot: VIA Technologies
- Dictionary: sxc.hu/topfer
- C870 GPU: Nvidia Corp.
- Old Books: flickr.com/ppdigital (cc)
- OpenCL Logo: Apple Corp./Ars Technica
- OS Platforms: flickr.com/aOliN.Tk
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- Queue: sxc.hu/cobrasoft
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- AMD logo: AMD Corporation
- Intel logo: Intel Corporation
- Cluster: sxc.hu/svilen001